

**DTU**





**DTU Health Technology  
Bioinformatics**

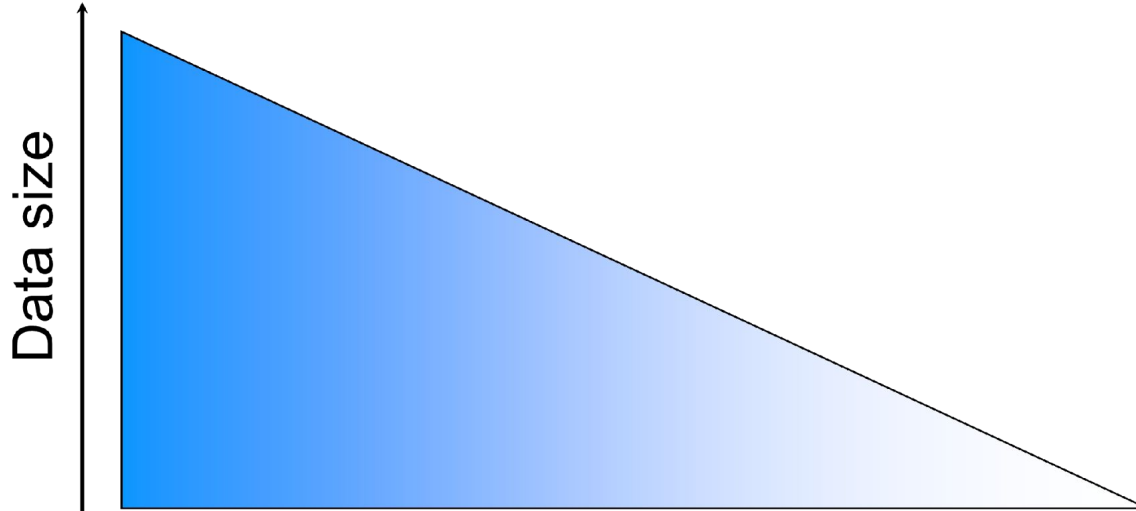
## **Alignment**

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# Menu

- Alignment approaches
- Burrows-Wheeler Transform
- More about coverage and depth
- Storing sequence alignments

# Generalized NGS analysis



Question

Raw  
reads

Pre-  
processing

Assembly:  
Alignment /  
*de novo*

Application  
specific:  
Variant calling,  
count matrix, ...

Compare  
samples /  
methods

Answer?



# What is an alignment?

Alignment = story

seq1 : CAAGACTAACCTGAA  
seq2 : CATGATAGCACTGCA

events?

seq1

CAAGACTAACCTGAA

seq2

CATGATAGCACTGCA



# What is an alignment?

Alignment = story

seq1 : CAAGACTAACCTGAA  
seq2 : CATGATAGCACTGCA

6 events

seq1

CAAGACTAACCTGAA

CATGACTAACCTGAA

CATGA TAACCTGAA

CATGA TAGACCTGAA

CATGA TAGCACCTGAA

CATGA TAGCA CTGAA

CATGA TAGCA CTGCA

seq2

CATGATAGCACTGCA



# What is an alignment?

Alignment = story

seq1: CAAGACTAACCTGAA

seq2: CATGATAGCACTGCA

D		C	A	T	G	A	T	A	G	C	A	C	T	G	C	A
	0	-2	-4	-6	-8	-10	-12	-14	-16	-18	-20	-22	-24	-26	-28	-30
C	-2	1	-1	-3	-5	-7	-9	-11	-13	-15	-17	-19	-21	-23	-25	-27
A	-4	-1	2	0	-2	-4	-6	-8	-10	-12	-14	-16	-18	-20	-22	-24
A	-6	-3	0	1	-1	-1	-3	-5	-7	-9	-11	-13	-15	-17	-19	-21
G	-8	-5	-2	-1	2	0	-2	-4	-4	-6	-8	-10	-12	-14	-16	-18
A	-10	-7	-4	-3	0	3	1	-1	-3	-5	-5	-7	-9	-11	-13	-15
C	-12	-9	-6	-5	-2	1	2	0	-2	-2	-4	-4	-6	-8	-10	-12
T	-14	-11	-8	-5	-4	-1	2	1	-1	-3	-3	-5	-3	-5	-7	-9
A	-16	-13	-10	-7	-6	-3	0	3	1	-1	-2	-4	-5	-4	-6	-6
A	-18	-15	-12	-9	-8	-5	-2	1	2	0	0	-2	-4	-6	-5	-5
C	-20	-17	-14	-11	-10	-7	-4	-1	0	3	1	-1	-3	-5	-6	-6
C	-22	-19	-16	-13	-12	-9	-6	-3	-2	1	2	2	0	-2	-2	-4
T	-24	-21	-18	-15	-14	-11	-8	-5	-4	-1	0	1	3	1	-1	-3
G	-26	-23	-20	-17	-14	-13	-10	-7	-4	-3	-2	-1	1	4	2	0
A	-28	-25	-22	-19	-16	-13	-12	-9	-6	-5	-2	-3	-1	2	3	3
A	-30	-27	-24	-21	-18	-15	-14	-11	-8	-7	-4	-3	-3	0	1	4

Needleman-Wunsch: best story?

- Dynamic programming
- Find the best path in the matrix



# What is an alignment?

Alignment = story

seq1: CAAGACTAACCTGAA  
seq2: CATGATAGCACTGCA

5 events

seq1

CAAGACTAACCTGAA

CATGACTAACCTGAA

CATGA TAACCTGAA

CATGA TAGCCTGAA

CATGA TAGCACTGAA

CATGA TAGCACTGCA

seq2

CATGATAGCACTGCA



# What is an alignment?

Alignment = story

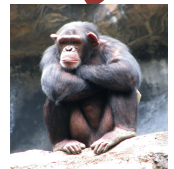
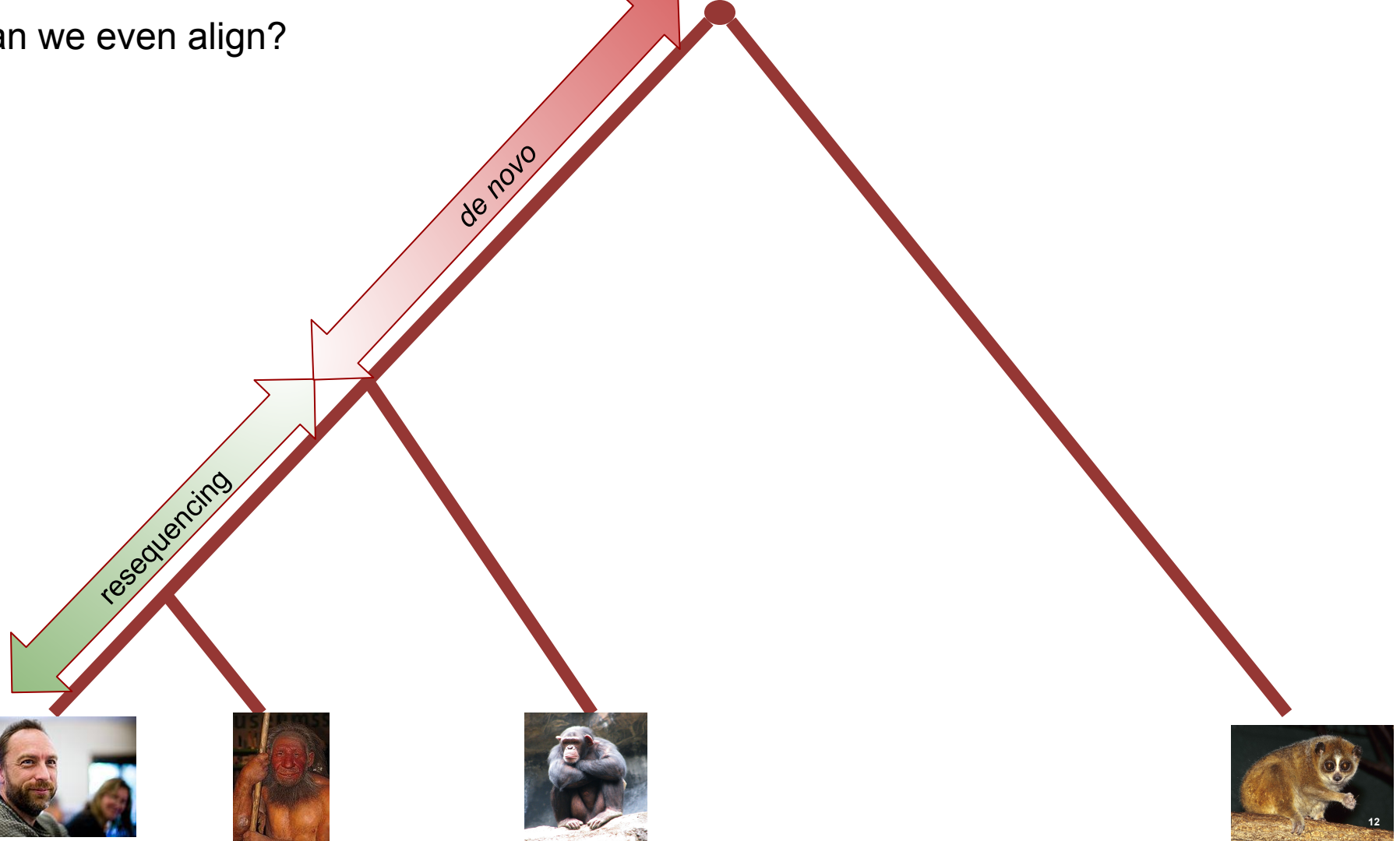
- 2 sequences can have a lot of alignments
- Not every alignment is equally likely
- Important to quantify the likelihood of seeing that alignment
- Be skeptical when you hear: “This is the alignment!”

# What is an alignment?

Types of alignment:

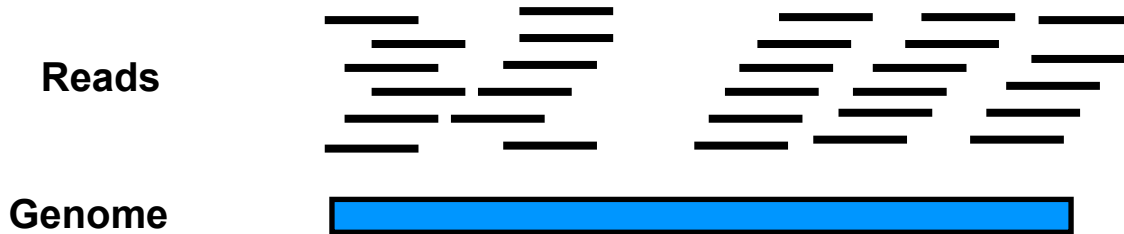
- “Short” read alignment
- Whole sequence alignment
- Whole genome/chromosome alignment
- Multiple sequence alignment

Can we even align?



# Alignment/Mapping

- Assemble your reads by aligning them to a closely related reference genome
- High sequence similarity between individuals makes this possible



## Sounds easy?

- Some pitfalls:
  - Divergence between sample and reference genome
  - Repeats in the genome
  - Recombination and re-arrangements
  - Poor reference genome quality
  - Read errors
  - Regions not in the ref. genome
  - Surprise sample



## Simplest solution

- Exact string matches!
- Calling “grep” on a read of 150bp on a 3B base genome takes only 1 second!
  - ... but aligning 20B reads from a NovaSeq would take 20B seconds so **634 years**
- We need bespoke tools
- What about mutations, sequencing errors etc?



Reference: ...ACGTGCGGACGCTGAACGTGA...

Read:                   | | | | | | | |  
                          GCGCACGCTGTAC

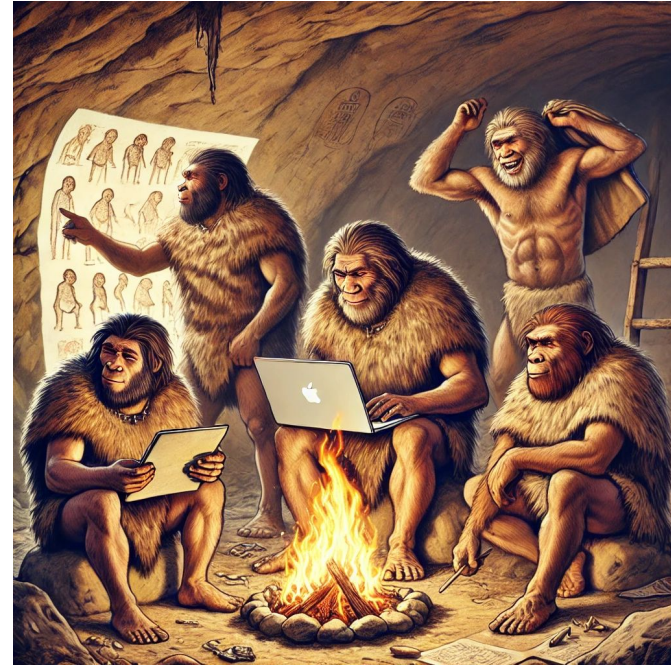
# A simpler solution?

- Allow imperfect matches

Reference: ...ACGTGCGGACGCTGAACGTGA...

||| ||||| ||  
Read: GCGCACGCTGTAC

- Use the Smith-Waterman algorithm to find the best match
- Aligning a read took 2.5 hours so aligning 20B reads from a NovaSeq would take **0.5M years**

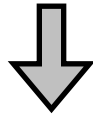




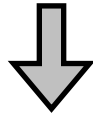


# Smart solution

1. Use algorithm to quickly find *possible* matches

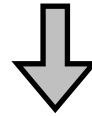


Drastically reduced search space

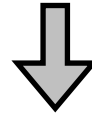


2. Allow us to perform slow/precise alignment for possible matches (Smith-Waterman)

3.2Gb



X possible matches



1 best match

# Hash based algorithms

Lookups in hashes are *fast!*

1. Index the reference using *k*-mers.
2. Search reads vs. hash *k*-mers
3. Perform alignment of entire read around seed
4. Report best alignment

Key	→	Value
.		.
.		.
.		.
ACTGCGTGTGA		Chr1_pos1234; Chr2_pos567
ACTGCGTGTGC		Chr7_posX
ACTGCGTGTGT		Chr7_posZ; ...
.		.
.		.
.		.

Also known as *Seed and extend*

## Spaced seeds

- Key/ $k$ -mer is called a seed
- Smarter: Spaced seeds (only care about '1' in seed, '0' = wildcards)
  - Higher sensitivity
  - One can use several seeds

111111111111

L = 11, 11 matches

111010010100110111

L = 18, 11 matches

## Multiple seeds & drawbacks

- One could require multiple short seeds
  - Instead of extending around each seed, extend around positions with several seed matches
- Drawbacks of hash-based approaches:
  - Lots(!) of RAM to keep index in memory (hg ~48Gb!)

# Burrows-Wheeler Transform

- Hash based aligners require lots of memory and are only reasonable fast
- Can we make it better/faster?
- Burrows Wheeler Transformation (BWT)
- BWT was originally created for compression

## BWT: Create index

Genome

Marks end-of-string, lexicographically smallest

T = AGGAGC\$

1. Create all possible shifts of the string

(move first base to end)

0 AGGAGC\$

## BWT: Create index

Genome

Marks end-of-string, lexicographically smallest

T = AGGAGC\$

1. Create all possible shifts of the string

(move first base to end)

0 AGGAGC\$

1 GGAGC\$A



## BWT: Create index

Genome

Marks end-of-string, lexicographically smallest

T = AGGAGC\$

1. Create all possible shifts of the string

(move first base to end)

0 AGGAGC\$  
1 GGAGC\$A  
2 GAGC\$AG

## BWT: Create index

Genome

Marks end-of-string, lexicographically smallest

T = AGGAGC\$

1. Create all possible shifts of the string

(move first base to end)

0	AGGAGC\$
1	GGAGC\$A
2	GAGC\$AG
3	AGC\$AGG

## BWT: Create index

Genome

Marks end-of-string, lexicographically smallest

T = AGGAGC\$

1. Create all possible shifts of the string

(move first base to end)

0	AGGAGC\$
1	GGAGC\$A
2	GAGC\$AG
3	AGC\$AGG
4	GC\$AGGA
5	C\$AGGAG
6	\$AGGAGC

# BWT: Create index

Genome

Marks end-of-string, lexicographically

smallest

T = AGGAGC\$

2. Sort the strings lexicographically to create BWT matrix and Suffix Array

0	AGGAGC\$
1	GGAGC\$A
2	GAGC\$AG
3	AGC\$AGG
4	GC\$AGGA
5	C\$AGGAG
6	\$AGGAGC

\$	A	G	G	A	G	C
A	G	C	\$	A	G	G
A	G	G	A	G	C	\$
C	\$	A	G	G	A	G
G	A	G	C	\$	A	G
G	C	\$	A	G	G	A
G	G	A	G	C	\$	A

BWT matrix

# BWT: Create index

Genome

Marks end-of-string, lexicographically smallest

T = AGGAGC\$

BWT (T) = CG\$GGAA

\$	A	G	G	A	G	C
A	G	C	\$	A	G	G
A	G	G	A	G	C	\$
C	\$	A	G	G	A	G
G	A	G	C	\$	A	G
G	C	\$	A	G	G	A
G	G	A	G	C	\$	A

BWT matrix

## BWT: Create index

Genome

Marks end-of-string, lexicographically

smallest

T = AGGAGC\$

- Reversible
- BTW(T) is easier to compress than T due to repeated characters tend to cluster ex:

BWT (T) = CG\$GGAA

Ringeren\_I\_Ringe\_ringer\_ringere\_end\_ringeren\_ringer\_i\_Ringsted\$  
\$d\_\_ \_nIiernerdenrgtrr\_gggggnnnnnnn\_RrrrRrReeeiiiiiiiieeeee\_\_\_\_gs

- try bzip2

# BWT: T-rank

T = AGGAGC\$

*T-ranking:*

*# of times the base  
occurred previously in T*

A<sub>0</sub> G<sub>0</sub> G<sub>1</sub> A<sub>1</sub> G<sub>2</sub> C<sub>0</sub> \$

<i>F</i>						<i>L</i>
\$	A	G	G	A	G	C
A	G	C	\$	A	G	G
A	G	G	A	G	C	\$
C	\$	A	G	G	A	G
G	A	G	C	\$	A	G
G	C	\$	A	G	G	A
G	G	A	G	C	\$	A

# BWT: *T*-rank

$T = \text{AGGAGC\$}$

*T*-ranking: # of times the base occurred previously in *T*

$A_0 \ G_0 \ G_1 \ A_1 \ G_2 \ C_0 \ \$$

<i>F</i>						<i>L</i>
\$	$A_0$	$G_0$	$G_1$	$A_1$	$G_2$	$C_0$
$A_1$	$G_2$	$C_0$	\$	$A_0$	$G_0$	$G_1$
$A_0$	$G_0$	$G_1$	$A_1$	$G_2$	$C_0$	\$
$C_0$	\$	$A_0$	$G_0$	$G_1$	$A_1$	$G_2$
$G_1$	$A_1$	$G_2$	$C_0$	\$	$A_0$	$G_0$
$G_2$	$C_0$	\$	$A_0$	$G_0$	$G_1$	$A_1$
$G_0$	$G_1$	$A_1$	$G_2$	$C_0$	\$	$A_0$



# BWT: T-rank

T = AGGAGC\$

*T-ranking: # of times the base occurred previously in T*

A<sub>0</sub> G<sub>0</sub> G<sub>1</sub> A<sub>1</sub> G<sub>2</sub> C<sub>0</sub> \$

<i>F</i>						<i>L</i>
\$	A <sub>0</sub>	G <sub>0</sub>	G <sub>1</sub>	A <sub>1</sub>	G <sub>2</sub>	C <sub>0</sub>
A <sub>1</sub>	G <sub>2</sub>	C <sub>0</sub>	\$	A <sub>0</sub>	G <sub>0</sub>	G <sub>1</sub>
A <sub>0</sub>	G <sub>0</sub>	G <sub>1</sub>	A <sub>1</sub>	G <sub>2</sub>	C <sub>0</sub>	\$
C <sub>0</sub>	\$	A <sub>0</sub>	G <sub>0</sub>	G <sub>1</sub>	A <sub>1</sub>	G <sub>2</sub>
G <sub>1</sub>	A <sub>1</sub>	G <sub>2</sub>	C <sub>0</sub>	\$	A <sub>0</sub>	G <sub>0</sub>
G <sub>2</sub>	C <sub>0</sub>	\$	A <sub>0</sub>	G <sub>0</sub>	G <sub>1</sub>	A <sub>1</sub>
G <sub>0</sub>	G <sub>1</sub>	A <sub>1</sub>	G <sub>2</sub>	C <sub>0</sub>	\$	A <sub>0</sub>

Notice that individual base-rank is the same in *F* and *L*



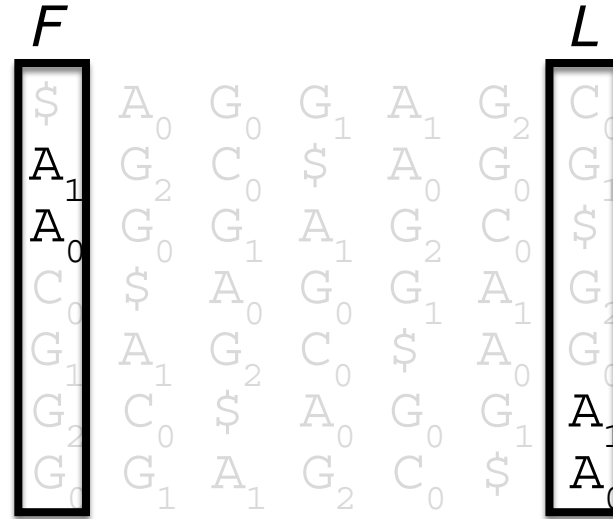
Rank will always be the same in *F* and *L*

# BWT: *T*-rank

$T = \text{AGGAGC\$}$

*T*-ranking: # of times the base occurred previously in *T*

$A_0 \ G_0 \ G_1 \ A_1 \ G_2 \ C_0 \ \$$



Notice that individual base-rank is the same in *F* and *L*



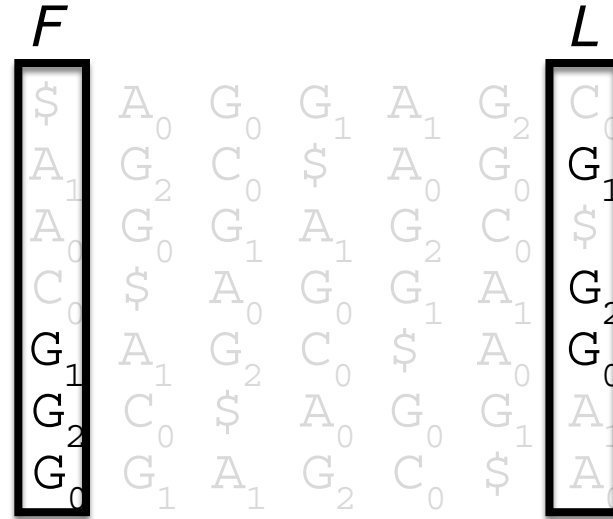
Rank will always be the same in *F* and *L*

# BWT: T-rank

T = AGGAGC\$

*T-ranking: # of times the base occurred previously in T*

A<sub>0</sub> G<sub>0</sub> G<sub>1</sub> A<sub>1</sub> G<sub>2</sub> C<sub>0</sub> \$



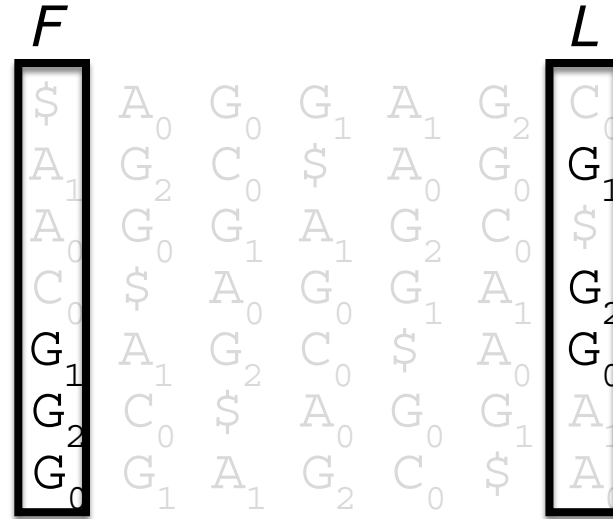
Notice that individual base-rank is the same in F and L



Rank will always be the same in F and L

# BWT: *T*-rank

Why does this generalize?



# BWT: *T*-rank

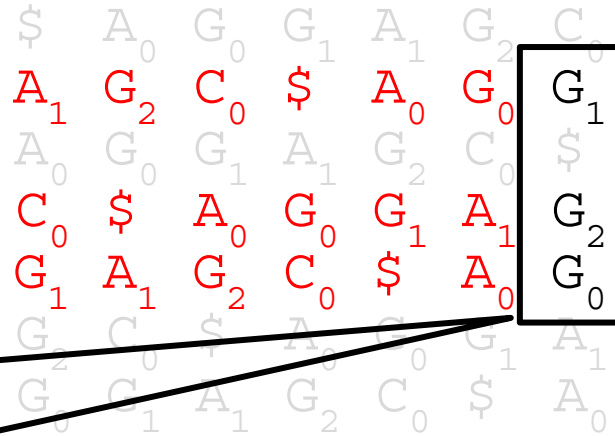
Why does this generalize?

We are all the same letter, our order is determined by what is right of us (blue).

<i>F</i>							<i>L</i>
\$	A <sub>0</sub>	G <sub>0</sub>	G <sub>1</sub>	A <sub>1</sub>	G <sub>2</sub>	C <sub>0</sub>	
A <sub>1</sub>	G <sub>2</sub>	C <sub>0</sub>	\$	A <sub>0</sub>	G <sub>0</sub>	G <sub>1</sub>	
A <sub>0</sub>	G <sub>0</sub>	G <sub>1</sub>	A <sub>1</sub>	G <sub>2</sub>	C <sub>0</sub>	\$	
C <sub>0</sub>	\$	A <sub>0</sub>	G <sub>0</sub>	G <sub>1</sub>	A <sub>1</sub>	G <sub>2</sub>	
G <sub>1</sub>	A <sub>1</sub>	G <sub>2</sub>	C <sub>0</sub>	\$	A <sub>0</sub>	G <sub>0</sub>	
G <sub>2</sub>	C <sub>0</sub>	\$	A <sub>0</sub>	G <sub>0</sub>	G <sub>1</sub>	A <sub>1</sub>	
G <sub>0</sub>	G <sub>1</sub>	A <sub>1</sub>	G <sub>2</sub>	C <sub>0</sub>	\$	A <sub>0</sub>	

# BWT: T-rank

Why does this generalize?



We are all the same letter, our order is determined by what is left of us (red).



## BWT: *T*-rank

How to reverse the BTW back the original string *T* ?

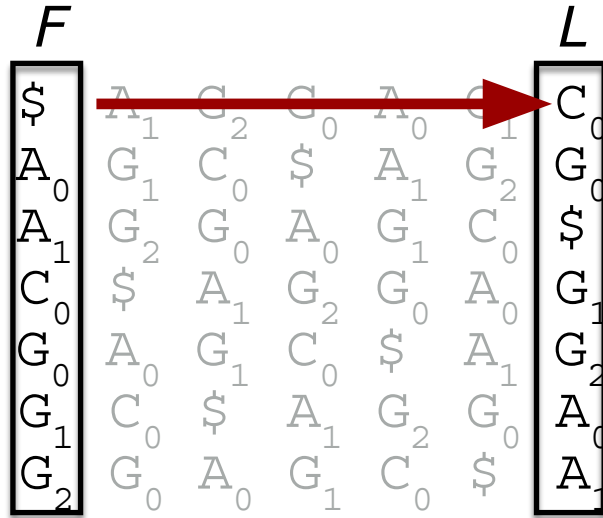
<i>F</i>	<i>T</i> -rank					<i>L</i>
\$	A <sub>1</sub>	G <sub>2</sub>	G <sub>0</sub>	A <sub>0</sub>	G <sub>1</sub>	C <sub>0</sub>
A <sub>0</sub>	G <sub>1</sub>	C <sub>0</sub>	\$	A <sub>1</sub>	G <sub>2</sub>	G <sub>0</sub>
A <sub>1</sub>	G <sub>2</sub>	G <sub>0</sub>	A <sub>0</sub>	G <sub>1</sub>	C <sub>0</sub>	\$
C <sub>0</sub>	\$	A <sub>1</sub>	G <sub>2</sub>	G <sub>0</sub>	A <sub>0</sub>	G <sub>1</sub>
G <sub>0</sub>	A <sub>0</sub>	G <sub>1</sub>	C <sub>0</sub>	\$	A <sub>1</sub>	G <sub>2</sub>
G <sub>1</sub>	C <sub>0</sub>	\$	A <sub>1</sub>	G <sub>2</sub>	G <sub>0</sub>	A <sub>0</sub>
G <sub>2</sub>	G <sub>0</sub>	A <sub>0</sub>	G <sub>1</sub>	C <sub>0</sub>	\$	A <sub>1</sub>



# BWT is reversible

LF-mapping: LF can be used to recreate the original genome

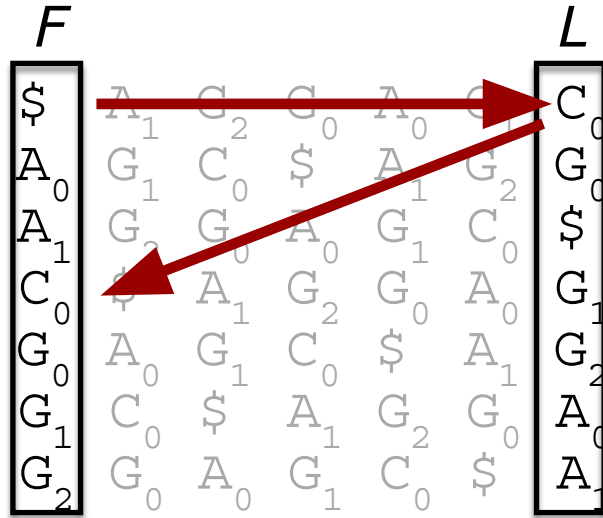
$C_0$



# BWT is reversible

LF-mapping: LF can be used to recreate the original genome

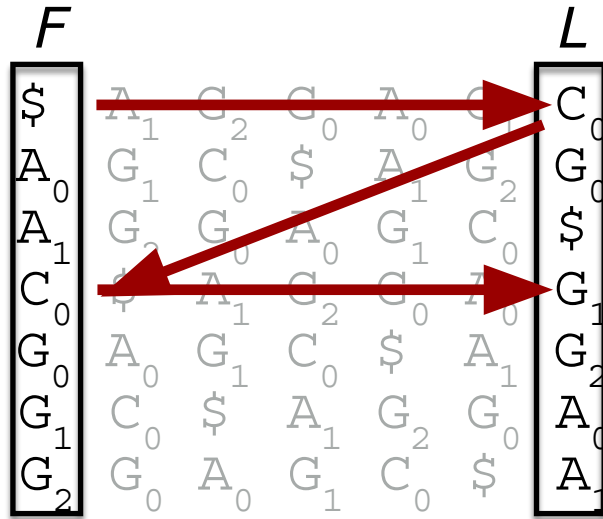
$C_0$



# BWT is reversible

LF-mapping: LF can be used to recreate the original genome

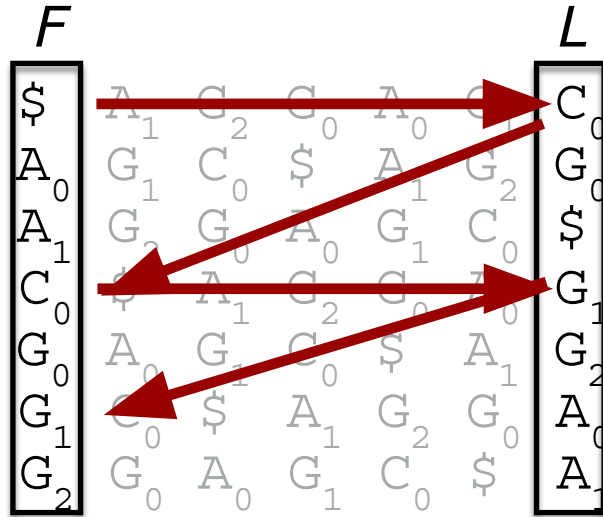
$C_0G_1$



# BWT is reversible

LF-mapping: LF can be used to recreate the original genome

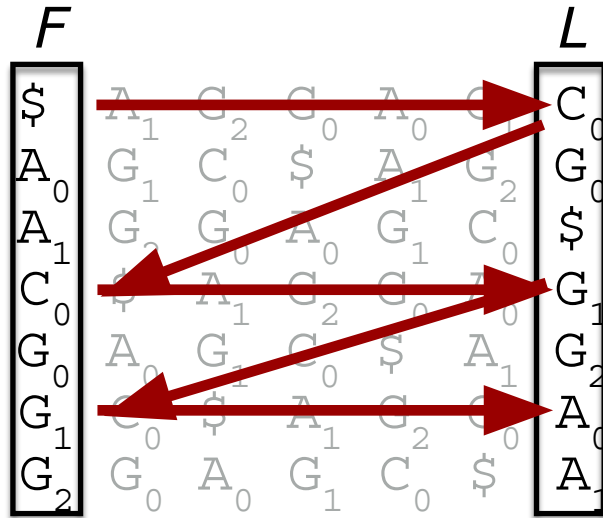
$C_0G_1$



# BWT is reversible

LF-mapping: LF can be used to recreate the original genome

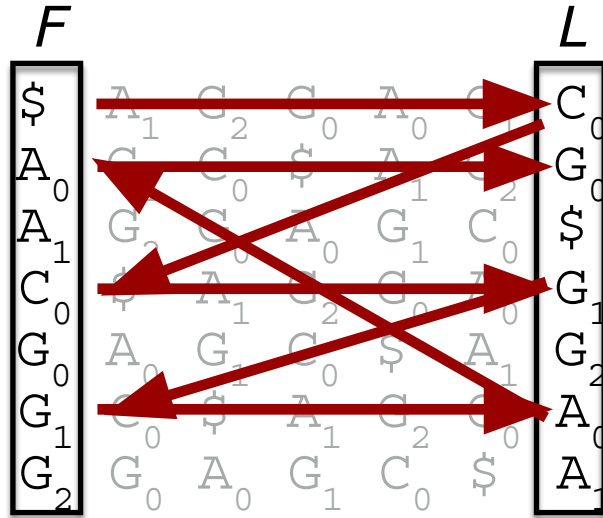
$C_0G_1A_0$



# BWT is reversible

LF-mapping: LF can be used to recreate the original genome

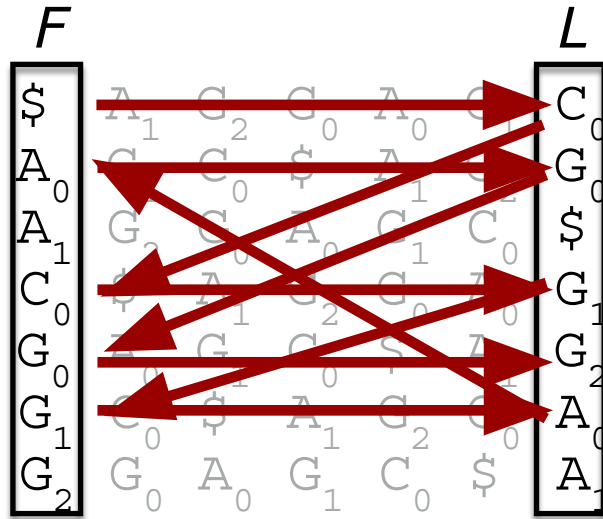
$C_0G_1A_0G_0$



# BWT is reversible

LF-mapping: LF can be used to recreate the original genome

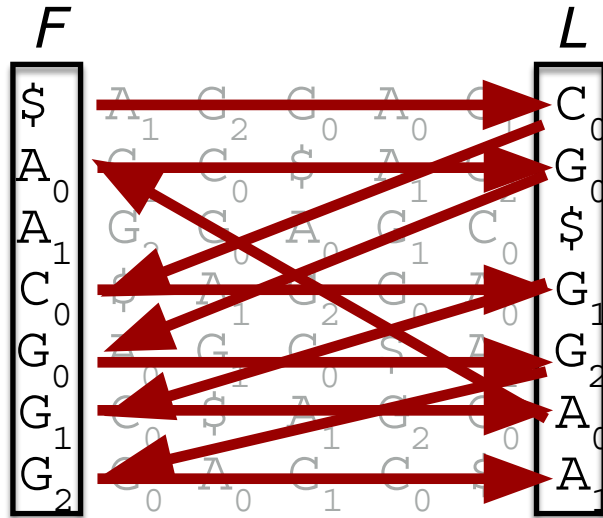
$C_0G_1A_0G_0G_2$



# BWT is reversible

LF-mapping: LF can be used to recreate the original genome

$C_0G_1A_0G_0G_2A_1$

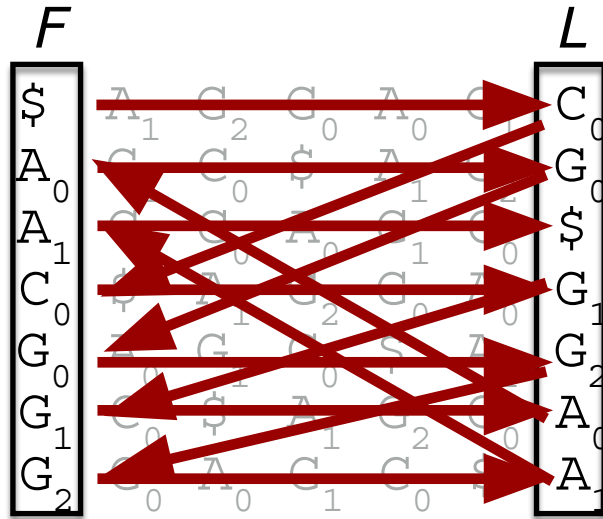




# BWT is reversible

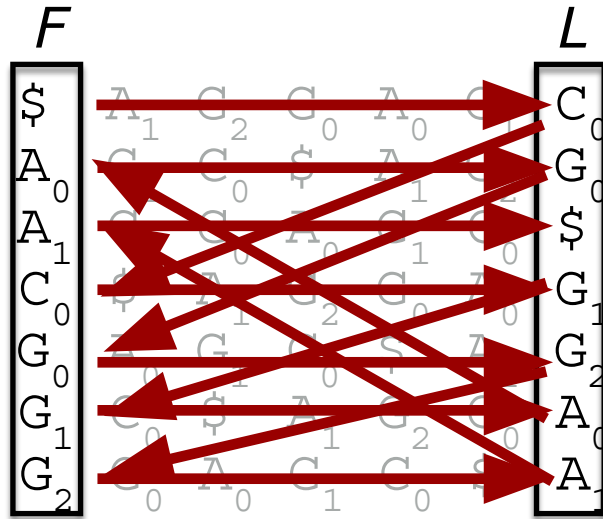
LF-mapping: LF can be used to recreate the original genome

$C_0G_1A_0G_0G_2A_1\$$



## BWT is reversible

LF-mapping: LF can be used to recreate the original genome



C<sub>0</sub>G<sub>1</sub>A<sub>0</sub>G<sub>0</sub>G<sub>2</sub>A<sub>1</sub>\$

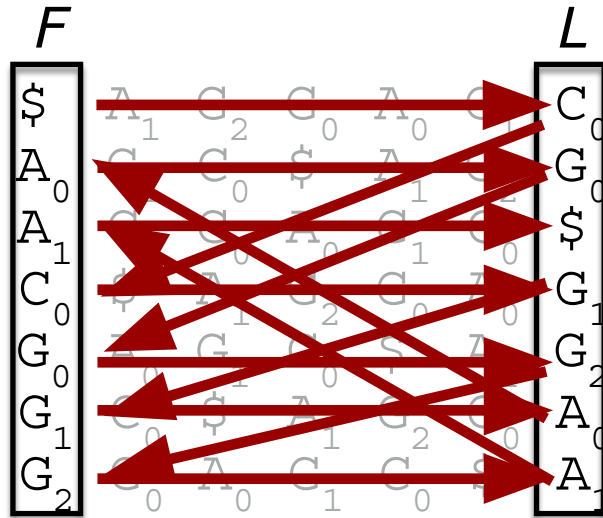
Reversed:

A<sub>1</sub>G<sub>2</sub>G<sub>0</sub>A<sub>0</sub>G<sub>1</sub>C<sub>0</sub>

T = AGGAGC\$

## BWT is reversible

LF-mapping: LF can be used to recreate the original genome



C<sub>0</sub>G<sub>1</sub>A<sub>0</sub>G<sub>0</sub>G<sub>2</sub>A<sub>1</sub>\$

Reversed:

A<sub>1</sub>G<sub>2</sub>G<sub>0</sub>A<sub>0</sub>G<sub>1</sub>C<sub>0</sub>

T = AGGAGC\$

F can be represented = 2x A, 1x C, 3x G  
we need  $|\Sigma|$  integers

We therefore only need to store  $L$

# Why are we talking about BTW for alignments?

- In 1994, Michael Burrows and David Wheeler created the Burrows-Wheeler Transform (BWT)
  - A reversible transformation of the genome
- Full-text index in Minute space (FM) index
  - *Paolo Ferragina, and Giovanni Manzini. "Opportunistic data structures with applications." Foundations of Computer Science, 2000. Proceedings. 41st Annual Symposium on. IEEE, 2000.*
  - Implementations: BWA, bowtie and SOAP2
  - How fast? First an intro to “big O” notation

## Brief intro to “Big O” notation

The CALCULATED procedure in Figure 3 gives a better, though not optimal, bound. It is conceptually equivalent to the one described in Figure 4, which is simpler to understand. We use the BWT of the reverse (not complemented) reference sequence to test if a substring of  $W$  is also a substring of  $X$ . Note that to do this test with BWT string  $B$  alone would make CALCULATED an  $O(|W|^2)$  procedure, rather than  $O(|W|)$  as is described in Figure 3.

?

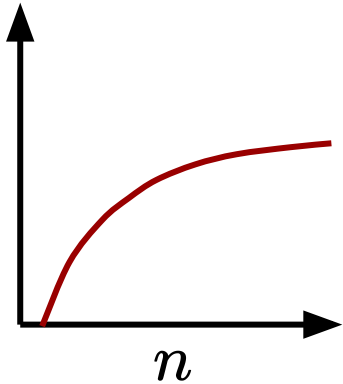


from: Li H, Durbin R. Fast and accurate short read alignment with Burrows-Wheeler transform. *Bioinformatics*. 2009 Jul 15;25(14):1754-60.

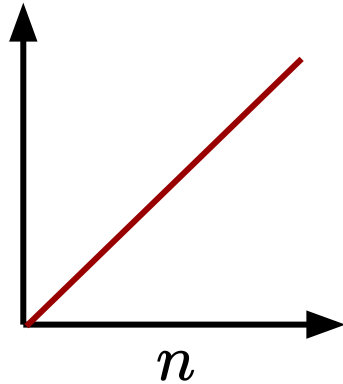
# Brief intro to “Big O” notation

- If I have  $n$  sequences, how does the amount of time required by the program increase?

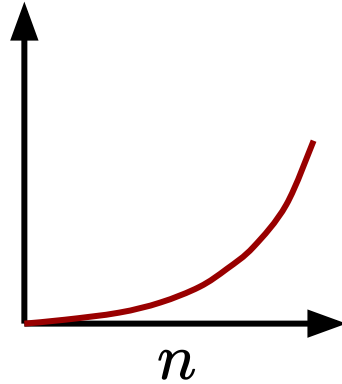
$\log(n)$



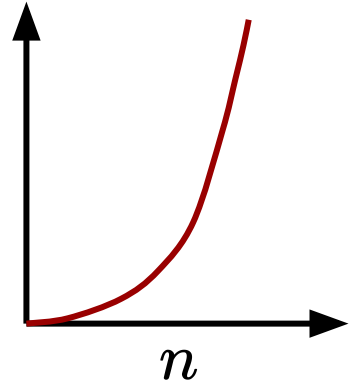
$n$



$n^2$



$2^n$



# Brief intro to “Big O” notation

- If I have  $n$  sequences, how does the amount of time required by the program increase?

Number of steps

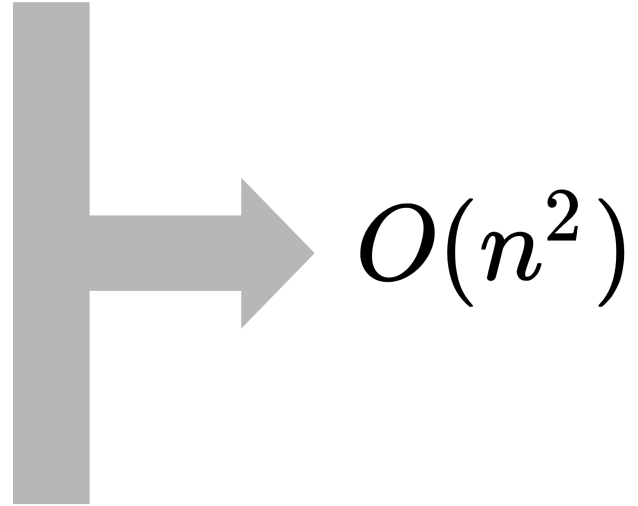
$$3n^2 + 6n + 1$$

$$25n^2 + 5n + 4$$

$$9n^2$$

$$1000n^2$$

$$n^2 - 9$$



$$O(n^2)$$

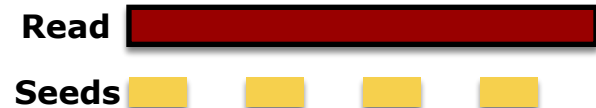
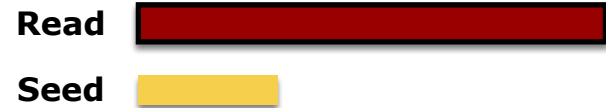
# BWT for alignment

- The FM index uses  $O(|L|)$  memory
- Can be searched in  $O(|pattern|)$  steps
- Entire FM-index is 1.5Gb for human genome
- FM-index
  - We also need certain other data structures that we did not cover
- Human genome can be effectively indexed and searched using 3Gb RAM!



# Implementation in BWA

- Burrows Wheeler Aligner (BWA) can use:
  - `bwa aln`: First ~30nt of read as seed
    - Extend around positions with seed match
    - For short reads
  - `bwa mem`: Multiple short seeds across the read
    - Extend around positions with several seed matches
    - For longer reads



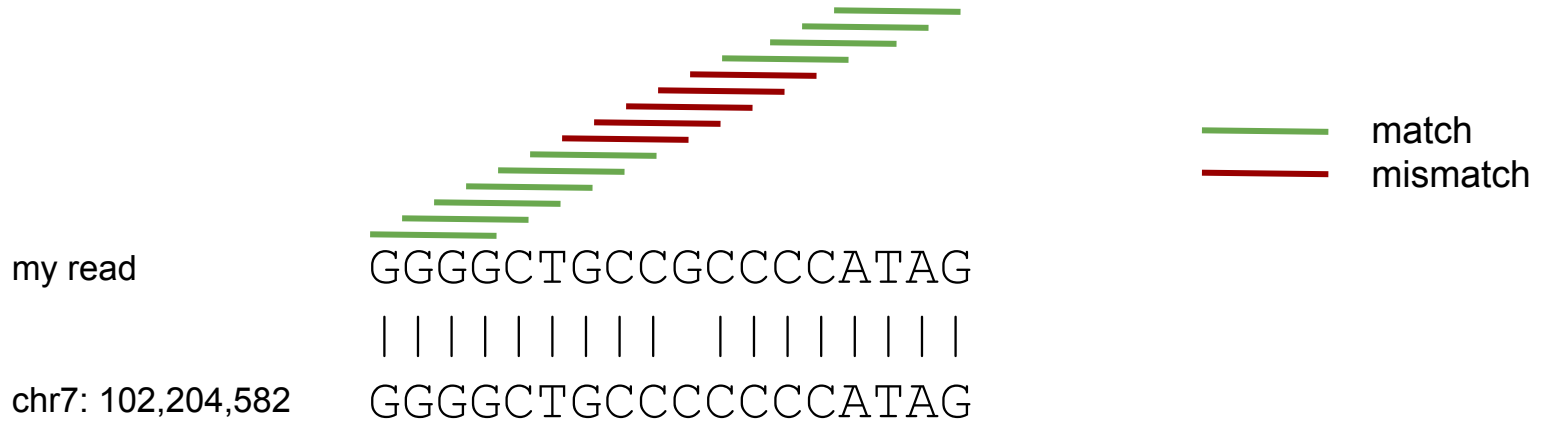
## Hash methods making a come back?

Let's go back to our example:

my read	GGGGCTGCCGCCCCATAG
chr7: 102,204,582	GGGGCTGCCGCCCCATAG

# Hash methods making a come back?

match with k=4



But wait,  $4^4$  is 256, that is too little!  
Everything will match, poor specificity

# Hash methods making a come back?

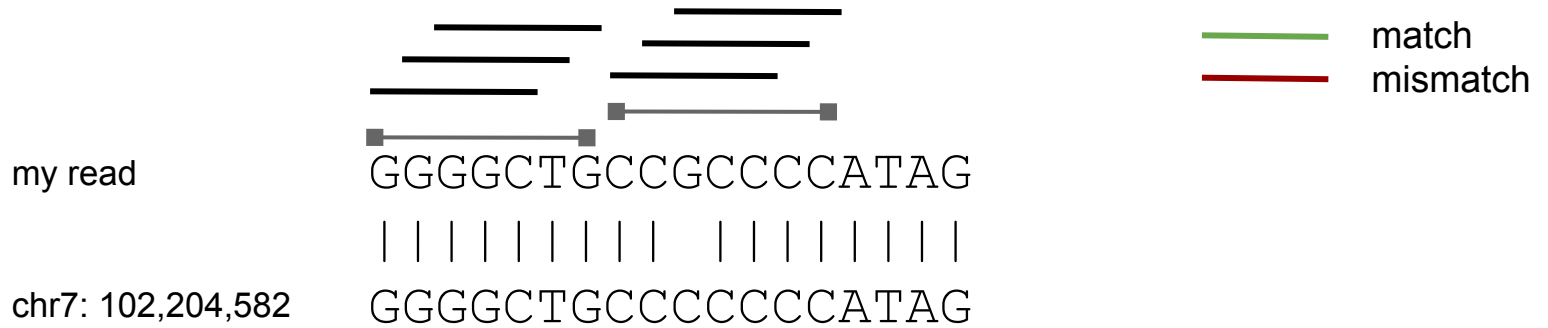
match with k=16



Nothing matches! poor sensitivity  
 $4^{16} = 4,294,967,296$

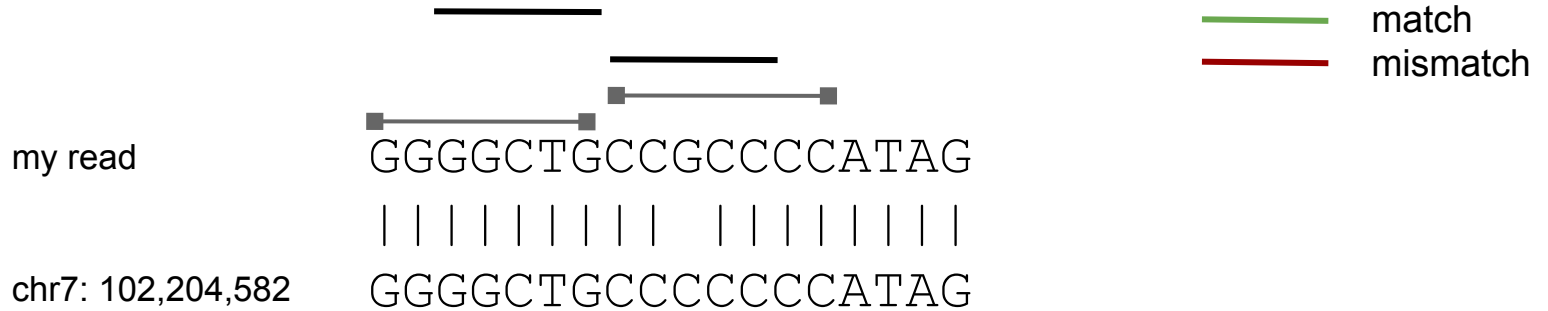
# Hash methods making a come back?

Mimimizer with  $w=7$   $k=5$



# Hash methods making a come back?

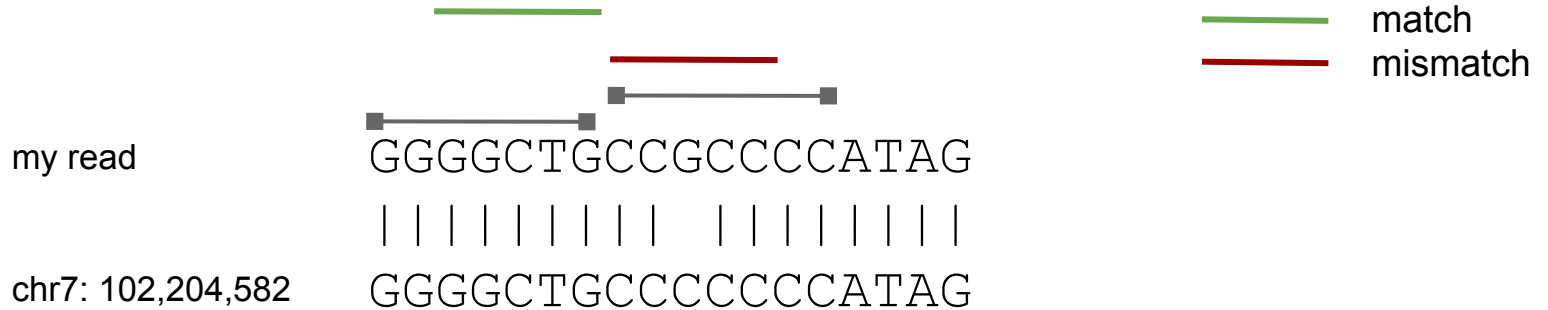
Mimimizer with  $w=7$   $k=5$



Pick the kmer with lowest lexicographical value

# Hash methods making a come back?

Mimimizer with  $w=7$   $k=5$



- Balance between short and long kmers
- Sensitive+specific

Great for long reads!

JOURNAL ARTICLE

## Minimap2: pairwise alignment for nucleotide sequences

Heng Li 

*Bioinformatics*, Volume 34, Issue 18, September 2018, Pages 3094–3100,

<https://doi.org/10.1093/bioinformatics/bty191>

**Published:** 10 May 2018 **Article history** ▼




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### Abstract

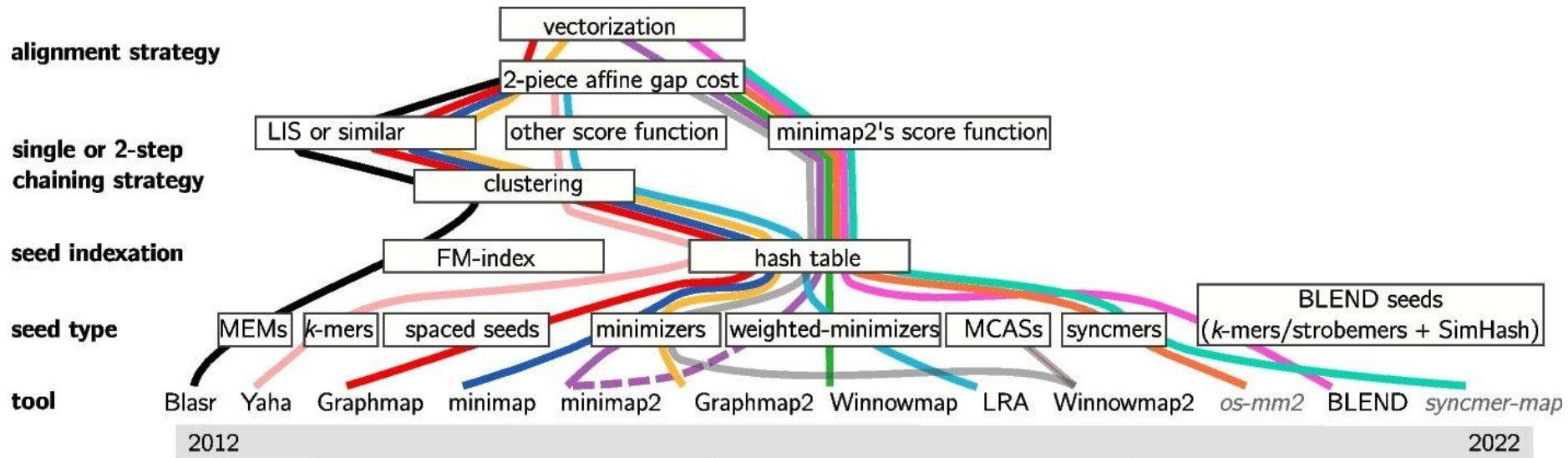
#### Motivation

Recent advances in sequencing technologies promise ultra-long reads of ~100 kb in average, full-length mRNA or cDNA reads in high throughput and genomic contigs over 100 Mb in length. Existing alignment programs are unable or inefficient to process such data at scale, which presses for the development of new alignment algorithms.

#### Results

Minimap2 is a general-purpose alignment program to map DNA or long mRNA sequences against a large reference database. It works with accurate short reads of  $\geq 100$  bp in length,  $\geq 1$  kb genomic reads at error rate ~15%, full-length noisy Direct RNA or cDNA reads and assembly contigs or closely related full chromosomes of hundreds of megabases in length. Minimap2 does split-read alignment, employs concave gap cost for long insertions and deletions and introduces new heuristics to reduce spurious alignments. It is 3–4 times as fast as mainstream short-read mappers at comparable accuracy, and is  $\geq 30$  times faster than long-read genomic or cDNA mappers at higher accuracy, surpassing most aligners specialized in one type of alignment.





Review | [Open access](#) | [Published: 01 June 2023](#)

## A survey of mapping algorithms in the long-reads era

[Kristoffer Sahlin](#) , [Thomas Baudeau](#), [Bastien Cazaux](#) & [Camille Marchet](#) 

*Genome Biology* 24, Article number: 133 (2023) | [Cite this article](#)

6594 Accesses | 1 Citations | 82 Altmetric | [Metrics](#)

### Abstract

It has been over a decade since the first publication of a method dedicated entirely to mapping long-reads. The distinctive characteristics of long reads resulted in methods moving from the seed-and-extend framework used for short reads to a seed-and-chain framework due to the seed abundance in each read. The main novelties are based on alternative seed constructs or chaining formulations. Dozens of tools now exist, whose heuristics have evolved considerably. We provide an overview of the methods used in long-read mappers. Since they are driven by implementation-specific parameters, we develop an original visualization tool to understand the parameter settings (<http://bczauz.polytech-lille.net/Minimap2/>).

## Intro to mapping quality

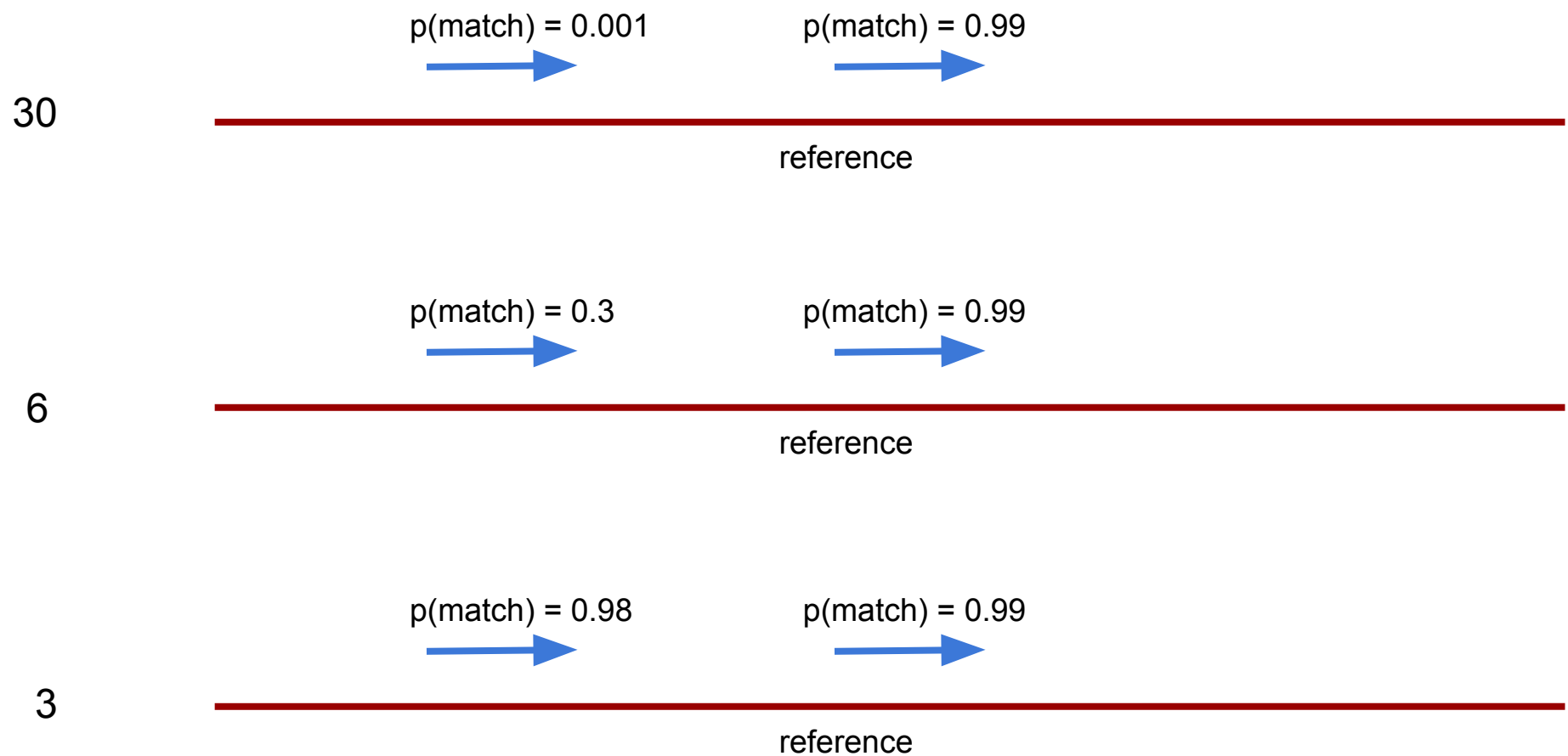
What happens when a sequence has multiple hits to the genome?  
Depends on the aligner, Burrows-Wheeler Aligner (BWA) does the following:

- Assign to the genomic location with the best score
- Use other matches to compute the probability of mismapping on a log scale:

$$\text{MAPQ} = -10 \log ( P[\text{mismapping}] )$$

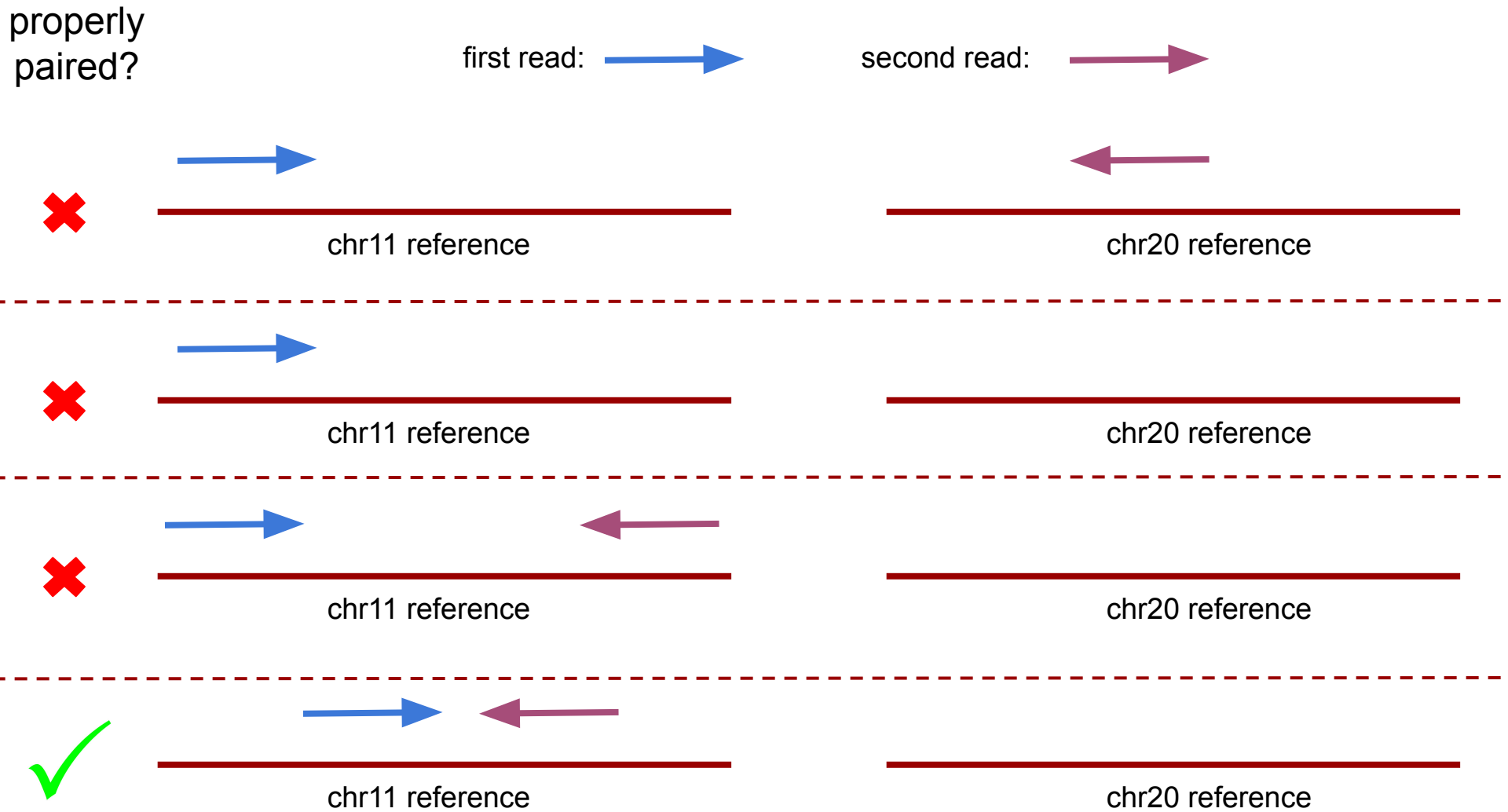
e.g. MAPQ 30 =  $P[\text{mismapping}] = 1/1000$

# Mapping quality



## Intro to “proper pairs” vs “unpaired”

- Some aligners add an extra flag to indicate that 2 paired reads were found:
  - on the same chromosome
  - facing each other (one + strand, the other - strand)
  - within a “reasonable” distance



## mapping quality vs mappability

- Mapping quality is often (poorly) approximated for speed
- Use another technique to avoid spurious mappings: genomic mappability
- Mapping quality is per read
- Mappability is for a genomic region

# Mappability

Ref. ATGCTGATGCTAGCGATATGCCCTAAAATCGATGCTAGCTGACTGATCGATCGACTGTCA

kmers of length 10

CCTAAAATCG	=1
CCCTAAAATC	=1
GCCCTAAAAT	=0.5
TGCCCTAAAA	=0.25
ATGCCCTAAA	=0.25
TATGCCCTAA	=0.5
ATATGCCCTA	=1
GATATGCCCT	=1
CGATATGCCC	=1
GCGATATGCC	=1



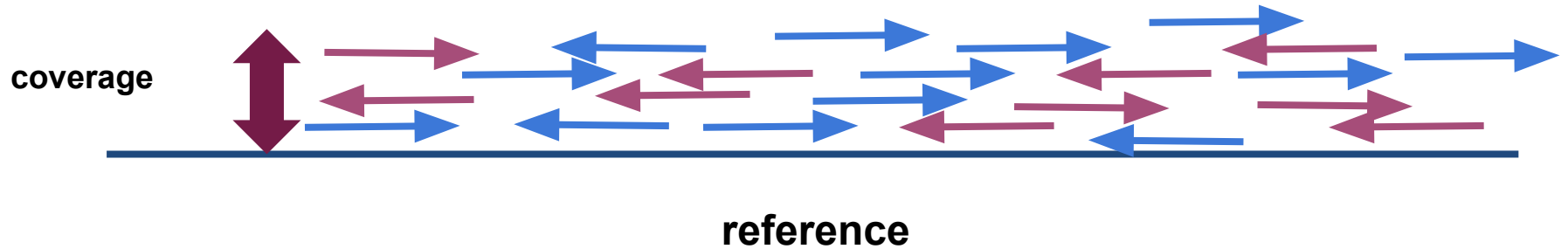
**Mean = 0.75**



mappability score 1 = unique <1=not unique

# Coverage

- Coverage/depth is how many times that your data covers the genome (on average)





# Depth of Coverage

reads:

G G C T A A

G C T A A G

G G G C T A A G G T

coverage

0 1 2 2 2 2 2 1 0 0

reference

# Depth of Coverage

average:  $(0+1+2+2+2+2+2+1+0+0)/10 = 1.2X$

reads:

G G C T A A

G C T A A G

G G G C T A A G G T

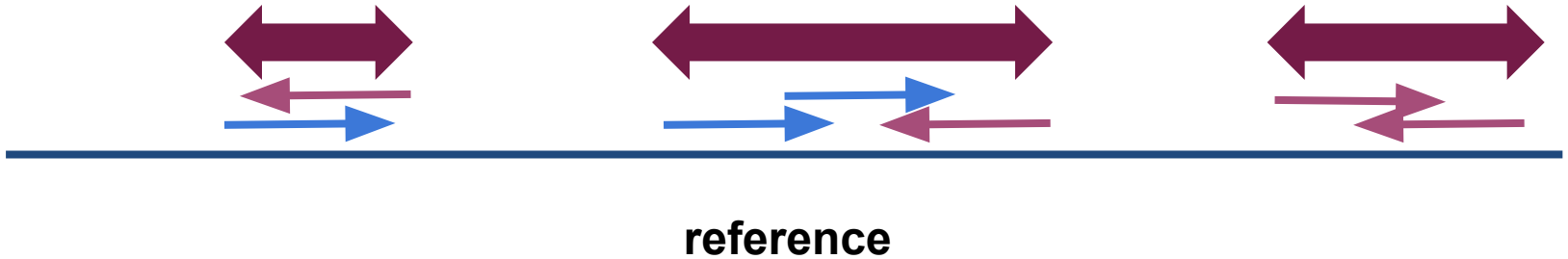
coverage

0 1 2 2 2 2 2 1 0 0

reference

# Breadth of coverage

- breadth of coverage: fraction of the genome covered (1X or more)



# Breadth of Coverage

average: 7 sites with Y / 10 = 70%

reads:

G G C T A A

G C T A A G

G G G C T A A G G T

covered?

N Y Y Y Y Y Y Y N N

reference

# Coverage

- Coverage/depth is how many times that your data covers the genome (on average)
- Example:
  - $N$ : Number of reads: 5M
  - $L$ : Read length: 100
  - $G$ : Genome size: 5Mbp
  - $C = 5M * 100 / 5M = 100X$
  - On average there are 100 reads covering each position in the genome

$$C = N \times \frac{L}{G}$$

# SAM/BAM format

- Sequence Alignment / Map format
- BAM = Binary SAM and zipped - always convert to BAM
- Two sections
  - Header: All lines start with “@”
  - Alignments: All other lines

```
@SQ SN:mito_ref LN:16569
@PG ID:bwa PN:bwa VN:0.7.17 CL:bwa samse human_MT.fa test_cut.sai input.fq.gz
HN3BX:8:2115:3985:4426 16 mito_ref 5265 37 12M * 0 0 ATTATCGAAGAA JJJJJAFFFAAA NM:i:0 MD:Z:12 RG:Z:sample1
XG:i:0
HN3BX:8:1102:2678:3157 0 mito_ref 9373 0 9M * 0 0 AATGATGAC AAFFFJJJJ NM:i:1 MD:Z:7G1 RG:Z:sample1
HN3BX:8:1217:1588:1335 16 mito_ref 7241 15 9M * 0 0 ATACACCAC JJJJFFFAA NM:i:0 MD:Z:9 RG:Z:sample2
XG:i:0
HN3BX:8:2216:6248:1342 16 mito_ref 14440 0 5M * 0 0 ATACT FFFAA NM:i:0 MD:Z:5 RG:Z:sample1
HN3BX:8:1222:8146:4237 0 mito_ref 1994 0 8M * 0 0 AACCTACC AAFFFJJJ NM:i:0 MD:Z:8 RG:Z:sample1
```

## The header section

```
@SQ SN:mito_ref LN:16569
@PG ID:bwa PN:bwa VN:0.7.17 CI:bwa samse human MT:fa test cut sai input fq.gz
HN3BX:8:2115:3985:4426 16 mito_ref 5265 37 12M * 0 0 ATTATCGAAGAA JJJJJJAFFFAAA NM:i:0 MD:Z:12 RG:Z:sample1
XG:i:0
HN3BX:8:1102:2678:3157 0 mito_ref 9373 0 9M * 0 0 AATGATGAC AAFFFJJJJ NM:i:1 MD:Z:7G1 RG:Z:sample1
HN3BX:8:1217:1588:1335 16 mito_ref 7241 15 9M * 0 0 ATACACCAC JJJJFFFAA NM:i:0 MD:Z:9 RG:Z:sample2
XG:i:0
HN3BX:8:2216:6248:1342 16 mito_ref 14440 0 5M * 0 0 ATACT FFFAA NM:i:0 MD:Z:5 RG:Z:sample1
HN3BX:8:1222:8146:4237 0 mito_ref 1994 0 5M * 0 0 AACCTAGC AAFFFJJJ NM:i:0 MD:Z:8 RG:Z:sample1
```

## The alignment section



```
@SQ SN:mito_ref LN:16569
PG ID:bwa PN:bwa VN:0.7.17 CL:bwa samse human_MT.fa test_cut.sai input fq.gz
```

Contains information like:

- What command line was used to generate this SAM/BAM file?
- What does the reference genomes look like? (chromosome names+length)

```
@SQ SN:mito_ref LN:16569
@PG ID:bwa PN:bwa VN:0.7.17 CL:bwa samse human_MT.fa test_cut.sai input.fq.gz
HN3BX:8:2115:3985:4426 16 mito_ref 5265 37 12M * 0 0 ATTATCGAAGAA JJJJJAFFFAAA NM:i:0 MD:Z:12 RG:Z:sample1
XG:i:0
HN3BX:8:1102:2678:3157 0 mito_ref 9373 0 9M * 0 0 AATGATGAC AAFFFJJJJ NM:i:1 MD:Z:7G1 RG:Z:sample1
HN3BX:8:1217:1588:1335 16 mito_ref 7241 15 9M * 0 0 ATACACCAC JJJJFFFAA NM:i:0 MD:Z:9 RG:Z:sample2
XG:i:0
HN3BX:8:2216:6248:1342 16 mito_ref 14440 0 5M * 0 0 ATACT FFFAA NM:i:0 MD:Z:5 RG:Z:sample1
HN3BX:8:1222:8146:4237 0 mito_ref 1994 0 8M * 0 0 AACCTACC AAFFFJJJ NM:i:0 MD:Z:8 RG:Z:sample1
```

HN3BX:8:2115:3985:442 XG:i:0	16	mito_ref	5265	37	12M	*	0	0	ATTATCGAAGAA	JJJJJJAFFFAAA	NM:i:0	MD:Z:12	RG:Z:sample1
HN3BX:8:1102:2678:315	0	mito_ref	9373	0	9M	*	0	0	AATGATGAC	AAFFFJJJJ	NM:i:1	MD:Z:7G1	RG:Z:sample1
HN3BX:8:1217:1588:133 XG:i:0	16	mito_ref	7241	15	9M	*	0	0	ATACACCAC	JJJJFFFAA	NM:i:0	MD:Z:9	RG:Z:sample2
HN3BX:8:2216:6248:134	16	mito_ref	14440	0	5M	*	0	0	ATACT	FFFAA	NM:i:0	MD:Z:5	RG:Z:sample1
HN3BX:8:1222:8146:4237	0	mito_ref	1994	0	8M	*	0	0	AACCTACC	AAFFFJJJJ	NM:i:0	MD:Z:8	RG:Z:sample1

Read IDs (should be unique)

```
HN3BX:8:2115:3985:4426 16 mito_ref 5265 37 12M * 0 0 ATTATCGAAGAA JJJJJAFFFFAAA NM:i:0 MD:Z:12 RG:Z:sample1
XG:i:0
HN3BX:8:1102:2678:3157 0 mito_ref 9373 0 9M * 0 0 AATGATGAC AAFFFJJJJ NM:i:1 MD:Z:7G1 RG:Z:sample1
HN3BX:8:1217:1588:1335 16 mito_ref 7241 15 9M * 0 0 ATACACCAC JJJJFFFAA NM:i:0 MD:Z:9 RG:Z:sample2
XG:i:0
HN3BX:8:2216:6248:1342 16 mito_ref 14440 0 5M * 0 0 ATACT FFFAA NM:i:0 MD:Z:5 RG:Z:sample1
HN3BX:8:1222:8146:4237 0 mito_ref 1994 0 8M * 0 0 AACCTACC AAFFFJJJ NM:i:0 MD:Z:8 RG:Z:sample1
```

flag (<https://broadinstitute.github.io/picard/explain-flags.html>)

ex:

0 = single-end read, mapped, read mapping to + strand  
16 = single-end read, mapped, read mapping to - strand

```
HN3BX:8:2115:3985:4426 15 mito_ref 5265 37 12M * 0 0 ATTATCGAAGAA JJJJJJAFFFAAA NM:i:0 MD:Z:12 RG:Z:sample1
XG:i:0
HN3BX:8:1102:2678:3157 0 mito_ref 9373 0 9M * 0 0 AATGATGAC AAFFFJJJJ NM:i:1 MD:Z:7G1 RG:Z:sample1
HN3BX:8:1217:1588:1335 15 mito_ref 7241 15 9M * 0 0 ATACACCAC JJJJFFFAA NM:i:0 MD:Z:9 RG:Z:sample2
XG:i:0
HN3BX:8:2216:6248:1342 15 mito_ref 14440 0 5M * 0 0 ATACT FFFAA NM:i:0 MD:Z:5 RG:Z:sample1
HN3BX:8:1222:8146:4237 0 mito_ref 1994 0 8M * 0 0 AACCTACC AAFFFJJJJ NM:i:0 MD:Z:8 RG:Z:sample1
```

name of chromosome (ex: chr1, chr2 ...)

```
HN3BX:8:2115:3985:4426 16 mito_re 5265 17 12M * 0 0 ATTATCGAAGAA JJJJJJAFFFAAA NM:i:0 MD:Z:12 RG:Z:sample1
XG:i:0
HN3BX:8:1102:2678:3157 0 mito_re 9373 10 9M * 0 0 AATGATGAC AAFFFJJJJ NM:i:1 MD:Z:7G1 RG:Z:sample1
HN3BX:8:1217:1588:1335 16 mito_re 7241 15 9M * 0 0 ATACACCAC JJJJFFFAA NM:i:0 MD:Z:9 RG:Z:sample2
XG:i:0
HN3BX:8:2216:6248:1342 16 mito_re 14440 10 5M * 0 0 ATACT FFFAA NM:i:0 MD:Z:5 RG:Z:sample1
HN3BX:8:1222:8146:4237 0 mito_re 1004 10 8M * 0 0 AACCTACC AAFFFJJJ NM:i:0 MD:Z:8 RG:Z:sample1
```

coordinate on chromosome

```
HN3BX:8:2115:3985:4426 16 mito_ref 5265 37 12M * 0 0 ATTATCGAAGAA JJJJJAFFFAAA NM:i:0 MD:Z:12 RG:Z:sample1
XG:i:0
HN3BX:8:1102:2678:3157 0 mito_ref 9373 0 9M * 0 0 AATGATGAC AAFFFJJJJ NM:i:1 MD:Z:7G1 RG:Z:sample1
HN3BX:8:1217:1588:1335 16 mito_ref 7241 15 9M * 0 0 ATACACCAC JJJJFFFAA NM:i:0 MD:Z:9 RG:Z:sample2
XG:i:0
HN3BX:8:2216:6248:1342 16 mito_ref 14440 0 5M * 0 0 ATACT FFFAA NM:i:0 MD:Z:5 RG:Z:sample1
HN3BX:8:1222:8146:4237 0 mito_ref 1994 0 3M * 0 0 AACCTACC AAFFFJJJ NM:i:0 MD:Z:8 RG:Z:sample1
```

mapping quality

ex:

MQ = prob of correct mapping

37 = 99.98%

15 = 96.84%

0 = 00.00%

```
HN3BX:8:2115:3985:4426 16 mito_ref 5265 37 12M * 0 0 ATTATCGAAGAA JJJJJAFFFAAA NM:i:0 MD:Z:12 RG:Z:sample1
XG:i:0
HN3BX:8:1102:2678:3157 0 mito_ref 9373 0 9M * 0 0 AATGATGAC AAFFJJJJJ NM:i:1 MD:Z:7G1 RG:Z:sample1
HN3BX:8:1217:1588:1335 16 mito_ref 7241 15 9M * 0 0 ATACACCAC JJJJFFFAA NM:i:0 MD:Z:9 RG:Z:sample2
XG:i:0
HN3BX:8:2216:6248:1342 16 mito_ref 14440 0 5M * 0 0 ATACT FFFAA NM:i:0 MD:Z:5 RG:Z:sample1
HN3BX:8:1222:8146:4237 0 mito_ref 1994 0 8M * 0 0 AACCTACC AAFFJJJJ NM:i:0 MD:Z:8 RG:Z:sample1
```

## CIGAR

Number of operations:

matches (**M**)

insertions (**I**)

deletions (**D**)



```
HN3BX:8:2115:3985:4426 16 mito_ref 5265 37 12 1 * 0 0 ATTATCGAAGAA JJJJJJAFFFAAA NM:i:0 MD:Z:12 RG:Z:sample1
XG:i:0
HN3BX:8:1102:2678:3157 0 mito_ref 9373 0 9 1 * 0 0 AATGATGAC AAFFFJJJJ NM:i:1 MD:Z:7G1 RG:Z:sample1
HN3BX:8:1217:1588:1335 16 mito_ref 7241 15 9 1 * 0 0 ATACACCAC JJJJFFFAA NM:i:0 MD:Z:9 RG:Z:sample2
XG:i:0
HN3BX:8:2216:6248:1342 16 mito_ref 14440 0 5 1 * 0 0 ATACT FFFAA NM:i:0 MD:Z:5 RG:Z:sample1
HN3BX:8:1222:8146:4237 0 mito_ref 1994 0 8 1 * 0 0 AACCTACC AAFFFJJJJ NM:i:0 MD:Z:8 RG:Z:sample1
```

chromosome of the other pair (not used here)

```
HN3BX:8:2115:3985:4426 16 mito_ref 5265 37 12M 0 ) ATTATCGAAGAA JJJJJAFFFAAA NM:i:0 MD:Z:12 RG:Z:sample1
XG:i:0
HN3BX:8:1102:2678:3157 0 mito_ref 9373 0 9M 0 ) AATGATGAC AAFFFJJJJ NM:i:1 MD:Z:7G1 RG:Z:sample1
HN3BX:8:1217:1588:1335 16 mito_ref 7241 15 9M 0 ) ATACACCAC JJJJFFFAA NM:i:0 MD:Z:9 RG:Z:sample2
XG:i:0
HN3BX:8:2216:6248:1342 16 mito_ref 14440 0 5M 0 ) ATACT FFFAA NM:i:0 MD:Z:5 RG:Z:sample1
HN3BX:8:1222:8146:4237 0 mito_ref 1994 0 8M 0 ) AACCTACC AAFFFJJJJ NM:i:0 MD:Z:8 RG:Z:sample1
```

position on the chromosome of the other pair (not used here)

```
HN3BX:8:2115:3985:4426 16 mito_ref 5265 37 12M * 0 TTATCGAAGAA JJJJJAFFFAAA NM:i:0 MD:Z:12 RG:Z:sample1
XG:i:0
HN3BX:8:1102:2678:3157 0 mito_ref 9373 0 9M * 0 ATGATGAC AAFFFJJJJ NM:i:1 MD:Z:7G1 RG:Z:sample1
HN3BX:8:1217:1588:1335 16 mito_ref 7241 15 9M * 0 TACACCAC JJJJFFFAA NM:i:0 MD:Z:9 RG:Z:sample2
XG:i:0
HN3BX:8:2216:6248:1342 16 mito_ref 14440 0 5M * 0 TACT FFFAA NM:i:0 MD:Z:5 RG:Z:sample1
HN3BX:8:1222:8146:4237 0 mito_ref 1994 0 8M * 0 0 AACCTACC AAFFFJJJJ NM:i:0 MD:Z:8 RG:Z:sample1
```

length between pairs (not used here)

```
HN3BX:8:2115:3985:4426 16 mito_ref 5265 37 12M * 0 0 ATTATCGAAGAA JJJJJJAFFFAAA NM:i:0 MD:Z:12 RG:Z:sample1
XG:i:0
HN3BX:8:1102:2678:3157 0 mito_ref 9373 0 9M * 0 0 AATGATGAC AAFFFJJJJ NM:i:1 MD:Z:7G1 RG:Z:sample1
HN3BX:8:1217:1588:1335 16 mito_ref 7241 15 9M * 0 0 ATACACCAC JJJJFFFAA NM:i:0 MD:Z:9 RG:Z:sample2
XG:i:0
HN3BX:8:2216:6248:1342 16 mito_ref 14440 0 5M * 0 0 ATACT FFFAA NM:i:0 MD:Z:5 RG:Z:sample1
HN3BX:8:1222:8146:4237 0 mito_ref 1994 0 8M * 0 0 AACCTACC AAFFFJJJ NM:i:0 MD:Z:8 RG:Z:sample1
```

sequence of the reads

```
HN3BX:8:2115:3985:4426 16 mito_ref 5265 37 12M * 0 0 ATTATCGAAGAA JJJJJAFFFAAA NM:i:0 MD:Z:12 RG:Z:sample1
XG:i:0
HN3BX:8:1102:2678:3157 0 mito_ref 9373 0 9M * 0 0 AATGATGAC AAFFFJJJJ NM:i:1 MD:Z:7G1 RG:Z:sample1
HN3BX:8:1217:1588:1335 16 mito_ref 7241 15 9M * 0 0 ATACACCAC JJJJFFFAA NM:i:0 MD:Z:9 RG:Z:sample2
XG:i:0
HN3BX:8:2216:6248:1342 16 mito_ref 14440 0 5M * 0 0 ATACT FFFAA NM:i:0 MD:Z:5 RG:Z:sample1
HN3BX:8:1222:8146:4237 0 mito_ref 1994 0 8M * 0 0 AACCTACC AAAAAIII NM:i:0 MD:Z:8 RG:Z:sample1
```

quality scores of the reads

```

HN3BX:8:2115:3985:4426 16 mito_ref 5265 37 12M * 0 0 ATTATCGAAGAA JJJJJJAFFFAAA NM:i:0 MD:Z:12 RG:Z:sample1
XG:i:0
HN3BX:8:1102:2678:3157 0 mito_ref 9373 0 9M * 0 0 AATGATGAC AAFFFJJJJ NM:i:1 MD:Z:7G1 RG:Z:sample1
HN3BX:8:1217:1588:1335 16 mito_ref 7241 15 9M * 0 0 ATACACCAC JJJJFFFAA NM:i:0 MD:Z:9 RG:Z:sample2
XG:i:0
HN3BX:8:2216:6248:1342 16 mito_ref 14440 0 5M * 0 0 ATACT FFFAA NM:i:0 MD:Z:5 RG:Z:sample1
HN3BX:8:1222:8146:4237 0 mito_ref 1994 0 8M * 0 0 AACCTACC AAFFFJJJ NM:i:0 MD:Z:8 RG:Z:sample1

```

optional flags, tell us:

- how many mismatches in alignment
- total number of matches
- read group

```
HN3BX:8:2115:3985:4426 16 mito_ref 5265 37 12M * 0 0 ATTATCGAAGAA JJJJJJAFFFFAAA NM:i:0 MD:Z:12 RG:Z:sample1
XG:i:0
HN3BX:8:1102:2678:3157 0 mito_ref 9373 0 9M * 0 0 AATGATGAC AAAAAJJJJJ NM:i:1 MD:Z:7G1 RG:Z:sample1
HN3BX:8:1217:1588:1335 16 mito_ref 7241 15 9M * 0 0 ATACACCAC JJJJFFFAA NM:i:0 MD:Z:9 RG:Z:sample2
XG:i:0
HN3BX:8:2216:6248:1342 16 mito_ref 14440 0 5M * 0 0 ATACT FFFAA NM:i:0 MD:Z:5 RG:Z:sample1
HN3BX:8:1222:8146:4237 0 mito_ref 1994 0 8M * 0 0 AACCTACC AAAAAJJJJ NM:i:0 MD:Z:8 RG:Z:sample1
```

**multiplexing: to which read group  
does the read belong to?**

# How to reconstruct the alignment? it's complicated...

```
HN3BX:8:2115:3985:4426 16 mito_ref 5265 37 12M * 0 0 ATTATCGAAGAA JJJJJAFFFAAA NM:i:0 MD:Z:12 RG:Z:sample1  
XG:i:0  
HN3BX:8:1102:2678:3157 0 mito_ref 9373 0 9M * 0 0 AATGATGAC AAFFFJJJJ NM:i:1 MD:Z:7G1 RG:Z:sample1  
HN3BX:8:1217:1588:1335 16 mito_ref 7241 15 9M * 0 0 ATACACCAC JJJJFFFAA NM:i:0 MD:Z:9 RG:Z:sample2  
XG:i:0  
HN3BX:8:2216:6248:1342 16 mito_ref 14440 0 5M * 0 0 ATACT FFFAA NM:i:0 MD:Z:5 RG:Z:sample1  
HN3BX:8:1222:8146:4237 0 mito_ref 1994 0 8M * 0 0 AACCTACC AAFFFJJJ NM:i:0 MD:Z:8 RG:Z:sample1
```

```
Ref+ :          9373 AATGATGGC 9381  
      | | | | | | | |  
Qry+ :          1  AATGATGAC  9
```



```
HN3BX:8:1102:2678:3157 0 mito_ref 9373 0 9M * 0 0 AATGATGAC AAFFFJJJJ NM:i:1 MD:Z:7G1 RG:Z:sample1
HN3BX:8:2115:3985:4426 16 mito_ref 5265 37 12M * 0 0 ATTATCGAAGAA JJJJJAFFFAAA NM:i:0 MD:Z:12 RG:Z:sample1
XG:i:0
HN3BX:8:1217:1588:1335 16 mito_ref 7241 15 9M * 0 0 ATACACCAC JJJJFFFAA NM:i:0 MD:Z:9 RG:Z:sample2
XG:i:0
HN3BX:8:2216:6248:1342 16 mito_ref 14440 0 5M * 0 0 ATACT FFFAA NM:i:0 MD:Z:5 RG:Z:sample1
HN3BX:8:1222:8146:4237 0 mito_ref 1994 0 AAFFFJJJ NM:i:0 MD:Z:8 RG:Z:sample1
```

sort by coordinate

```
HN3BX:8:1222:8146:4237 0 mito_ref 1994 0 8M * 0 0 AACCTACC AAFFFJJJ NM:i:0 MD:Z:8 RG:Z:sample1
HN3BX:8:2115:3985:4426 16 mito_ref 5265 7 12M * 0 0 ATTATCGAAGAA JJJJJAFFFAAA NM:i:0 MD:Z:12 RG:Z:sample1
XG:i:0
HN3BX:8:1217:1588:1335 16 mito_ref 7241 5 9M * 0 0 ATACACCAC JJJJFFFAA NM:i:0 MD:Z:9 RG:Z:sample2
XG:i:0
HN3BX:8:1102:2678:3157 0 mito_ref 9373 0 9M * 0 0 AATGATGAC AAFFFJJJJ NM:i:1 MD:Z:7G1 RG:Z:sample1
HN3BX:8:2216:6248:1342 16 mito_ref 14440 0 5M * 0 0 ATACT FFFAA NM:i:0 MD:Z:5 RG:Z:sample1
```

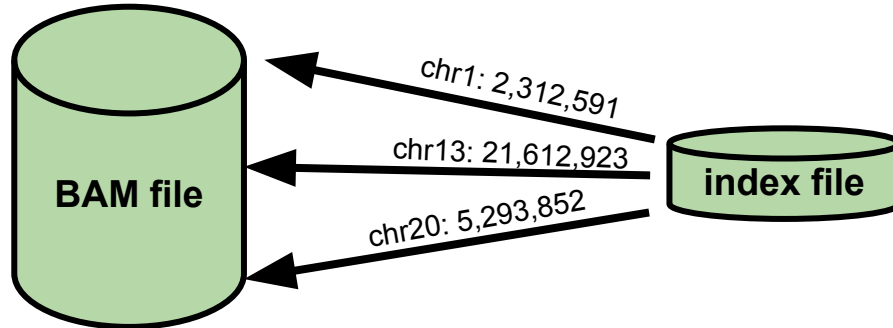
```
HN3BX:8:1102:2678:3157 0 mito_ref 9373 0 9M * 0 0 AATGATGAC AAAAAJJJJ NM:i:1 MD:Z:7G1 RG:Z:sample1
HN3BX:8:2115:3985:4426 16 mito_ref 5265 37 12M * 0 0 ATTATCGAAGAA JJJJJAAAA NM:i:0 MD:Z:12 RG:Z:sample1
XG:i:0
HN3BX:8:1217:1588:1335 16 mito_ref 7241 15 9M * 0 0 ATACACCAC JJJJJFFAA NM:i:0 MD:Z:9 RG:Z:sample2
XG:i:0
HN3BX:8:2216:6248:1342 16 mito_ref 14440 0 5M * 0 0 ATACT FFFAA NM:i:0 MD:Z:5 RG:Z:sample1
HN3BX:8:1222:8146:4237 0 mito_ref 1994 0 AAAAAJJJJ NM:i:0 MD:Z:8 RG:Z:sample1
```

sort by read name

```
HN3BX:8:1102:2678:3157 0 mito_ref 9373 0 9M * 0 0 AATGATGAC AAAAAJJJJ NM:i:1 MD:Z:7G1 RG:Z:sample1
HN3BX:8:1217:1588:1335 16 mito_ref 7241 15 9M * 0 0 ATACACCAC JJJJJFFAA NM:i:0 MD:Z:9 RG:Z:sample2
XG:i:0
HN3BX:8:1222:8146:4237 0 mito_ref 1994 0 8M * 0 0 AACCTACC AAAAAJJJJ NM:i:0 MD:Z:8 RG:Z:sample1
HN3BX:8:2115:3985:4426 16 mito_ref 5265 37 12M * 0 0 ATTATCGAAGAA JJJJJAAAA NM:i:0 MD:Z:12 RG:Z:sample1
XG:i:0
HN3BX:8:2216:6248:1342 16 mito_ref 14440 0 5M * 0 0 ATACT FFFAA NM:i:0 MD:Z:5 RG:Z:sample1
```

# BAM indexing

- I want all reads mapping to chrX
- Go through the entire BAM file until we reach chrX
- Better idea: keep a small file telling you where chromosome/coord start:



- Need a BAM file sorted by coordinate
- Is usually a .bai or .csi file

# BAM vs CRAM

- Idea: why store the original reads? You only the differences to a reference
- Can be 30-60% smaller
- Good for long term storage
- BAM can be converted to CRAM
- Can be indexed

## **Brief note about pangenomes graphs**

TCTGTAAT

CAACCTCA

CTGTTCTG

CCTCACCA

---

AGCCTGTTCTGTAATCGATAAACCCCGATCAACCTCACCA

reference

unmapped:

GACAAGTC



AAGTCCCG





TCTGTAAT

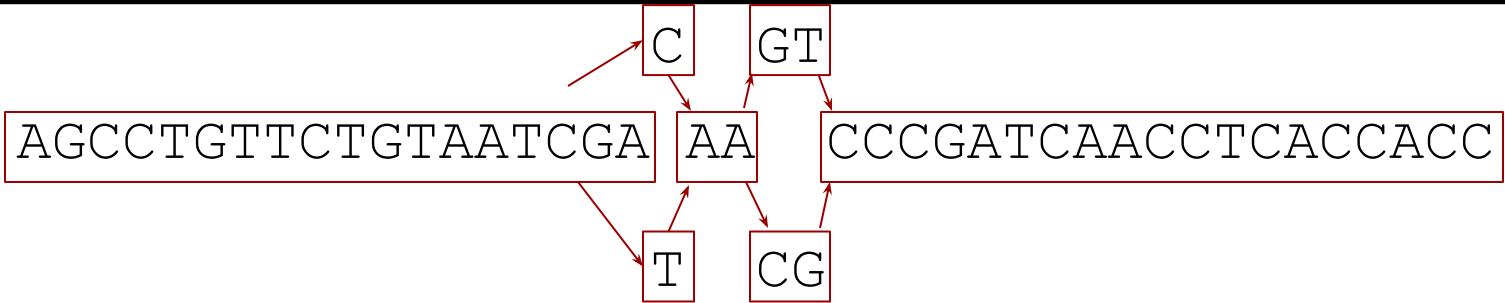
AAGTCCCG

CAACCTCA

CTGTACTG

GACAAGTC

CCTCACCA



references

unmapped:





TCTGTAAT

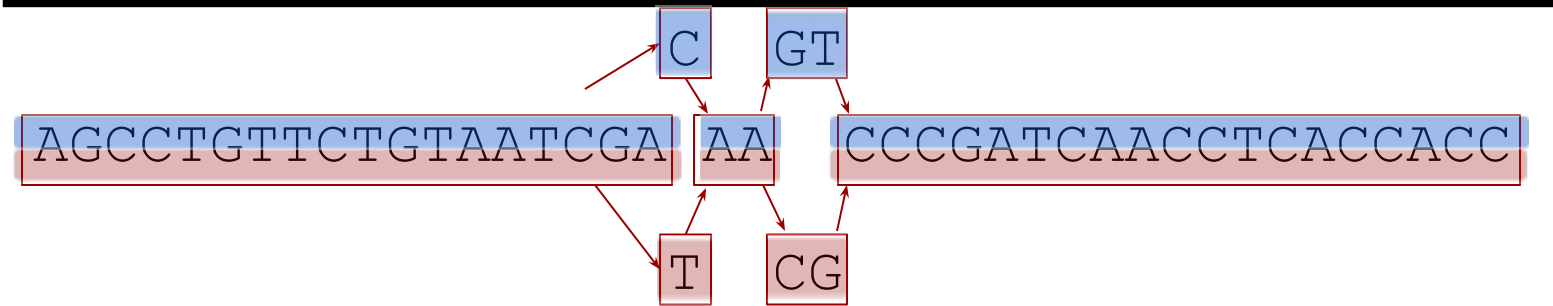
AAGTCCCG

CAACCTCA

CTGTACTG

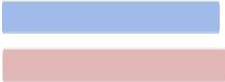
GACAAGTC

CCTCACCA



references

Paths:



unmapped:



**Exercise time!**