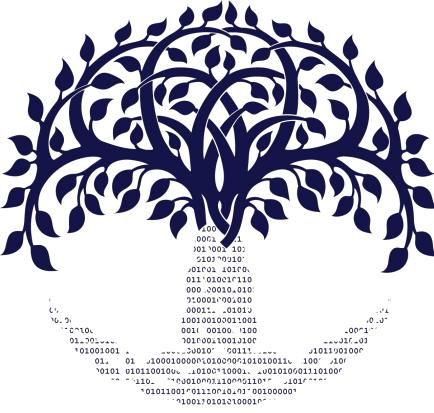


DTU





**DTU Health Technology
Bioinformatics**

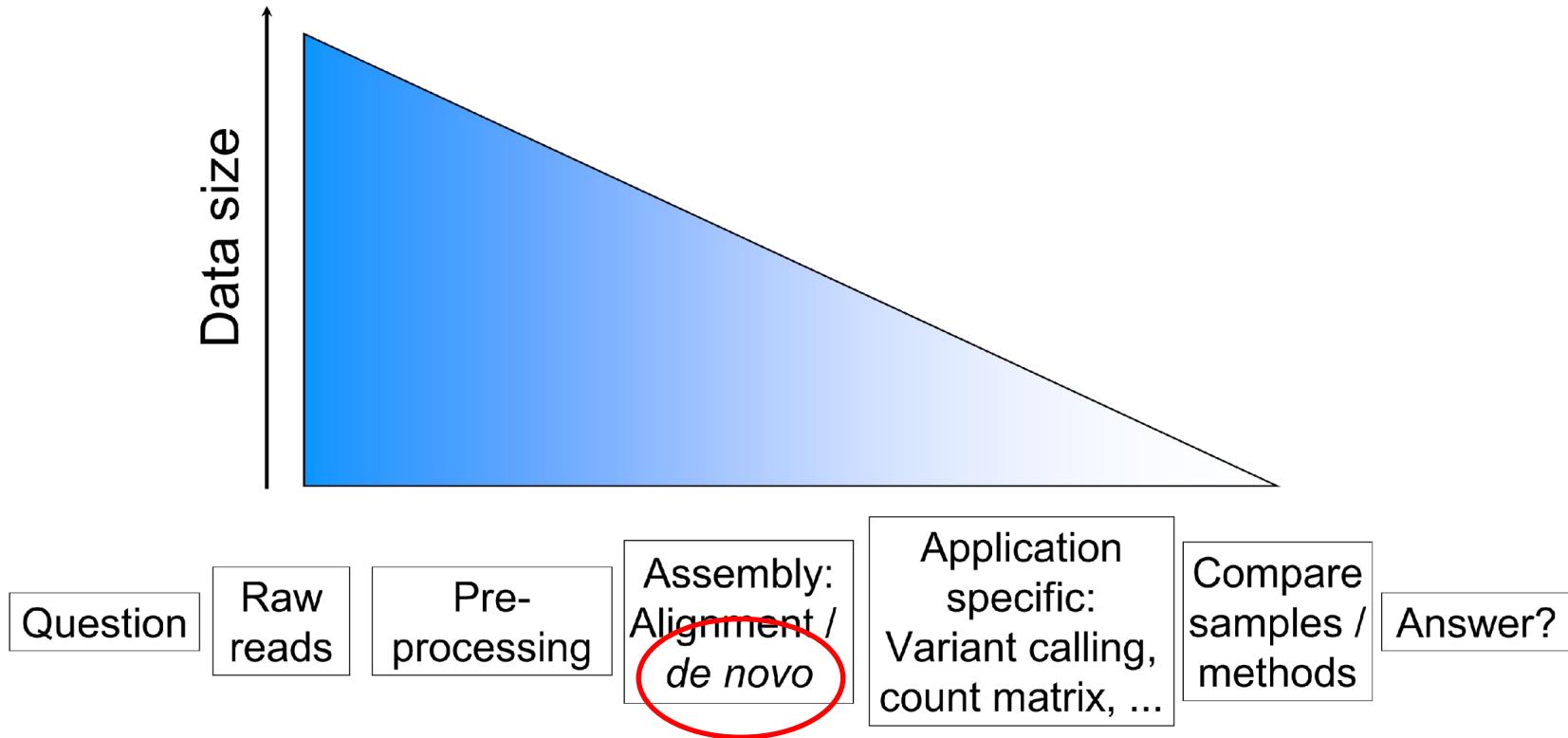
de novo assembly

*Gabriel Renaud
Associate Professor
Section of Bioinformatics
Technical University of Denmark
gisves@dtu.dk*

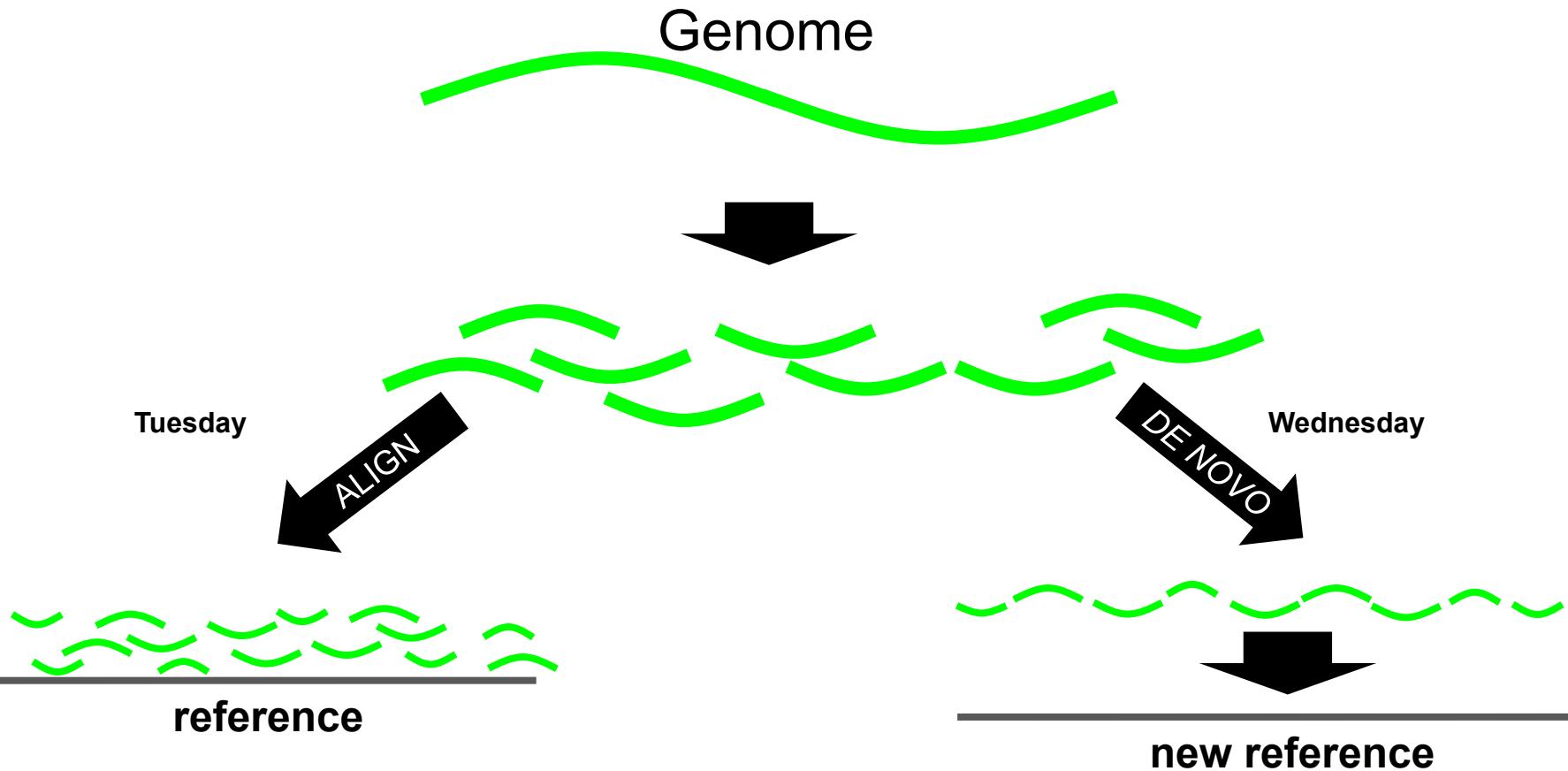
Menu

- Assembly approaches
- Assembly graphs
- Graph postprocessing filtering
- The woes of repetition
- Benchmarking your assembly

Generalized NGS analysis



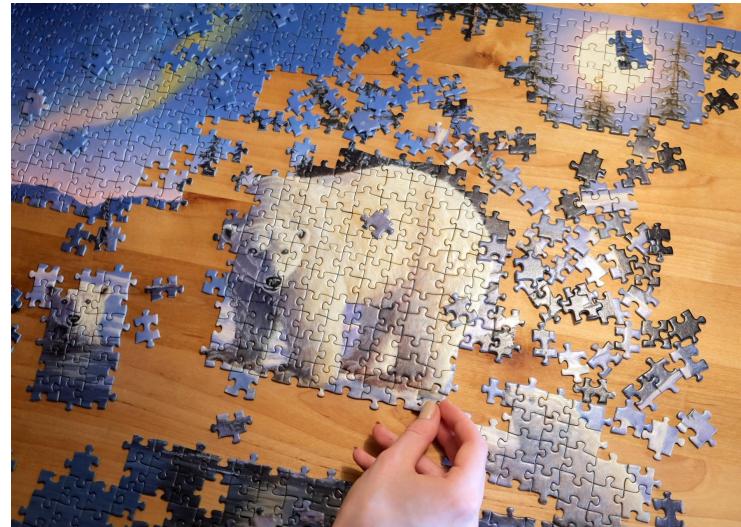
Whole genome sequencing



Input



Output



Input

```
@MISEQ423_0:+:7218:7278:60-2  
GTTACTCGGACTACCCGATGCATACACCACATGAAACA  
T  
+  
]V[P]]\]]]]]\]]]]]]]]]]\]]]\]]]]]]]]]]]]\]  
]  
@MISEQ423_0:-:15245:15305:60-2  
AGGGCAAGATGAAGTGAAGGTAAAGAATCGTGTGAGGG  
T  
+  
]]]][[Z]]]]]]]]]]]]]]]]]]]]]]]]]]]]]]]]]]]  
]  
@MISEQ423_0:-:242:302:60-2  
TTTGGTGGAAATTGGTTATGATGTCTGTGTGGAAAG  
T  
+  
]]]]]]]]Z]]]]]]]]]]]]Z]]]\]]Z]]]]]]]  
]  
@MISEQ423_0:-:1729:1789:60-2  
TGCAGGTACTATCTATTGCGCCAGGTTCAATTCTAT  
C  
+  
11111111X11111111111111111111111111111111
```



Output

```
>contig#25_0  
GATCACAGGTCTATCACCTATTAAACCACTCACGGGAGCTCTCCA  
GTATGCACGCGATAGCATTGCGAGACGCTGGAGCCGGAGCACCC  
CTGCCTCATCCTATTATTTATCGCACCTACGTTCAATATTACAGG  
ATTAATTAATGCTTAGGACATAATAACAATTGAATGTCTG  
ATAACAAAAAAATTCCACCAAACCCCCCTCCCCCGCTCTGGCC  
AACCCCCAAAACAAAGAACCTAACACACCAGCCTAACAGATTCA  
TTTAACAGTCACCCCCCAACTAACACATTATTTCCCCTCCCAC  
CAACCCCCGCCCATCCTACCCAGCACACACACCCGCTGCTAAC  
AAAGACACCCCCCACAGTTATGCTTACCTCCTCAAAGCAAT  
ACATCACCCCATAAACAAATAGGTTGGTCTAGCCTTCTATT  
GCATCCCCGTTCCAGTGAGTTCACCTCTAAATCACCACGATCAA  
AATGCAGCTAAAACGCTTAGCCTAGCCACACCCCCACGGGAAAC  
ACGAAAGTTAACTAAGCTATAACTAACCCCCAGGGT
```

Important definitions

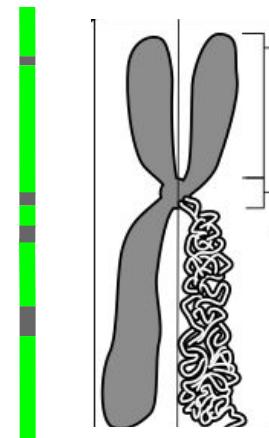
Contigs



Scaffolds



Chromosome



Important definitions

Contigs



```
>contig#1
GATCACAGGTCTATCACCTATTAAACCACTCACGGGAGCTCTCCA
GTATGCACGCGATAGCATTGCGAGACGCTGGAGCCGGAGCACCCCT
CTGCCTCATCCTATTATTATCGCACCTACGTTCAATATTACAGG
>contig#2
ATTAATTAATGCTTGTAGGACATAATAACAATTGAATGTCTG
ATAACAAAAAAATTCCACCAAACCCCCCTCCCCCGCTTCTGGCC
>contig#3
AACCCCCAAAACAAAGAACCTAACACCAGCCTAACCAAGATTCA
TTTTAACAGTCACCCCCCAACTAACACATTATTTCCCCTCCCAC
CAACCCCCGCCATCCTACCCAGCACACACACACCCTGCTGCTAAC
AAAGACACCCCCCACAGTTATGTAGCTTACCTCCTCAAAGCAAT
>contig#4
ACATCACCCCCATAAACAAATAGGTTGGTCTAGCCTTCTATTA
GCATCCCCGTTCCAGTGAGTTCACCTCTAAATCACCACGATCAA
AATGCAGCTAAAACGCTTAGCCTAGCCACACCCCCACGGGAAAC
ACGAAAGTTAACTAAGCTATACTAACCCAGGGT
```

Important definitions

```
>scaffold#1
AACCCCAAAACAAAGAACCTAACACCAGCCTAACAGAGATTCA
TTTTAACAGTCACCCCCCAACTAACACATTATTTCCCCTCCCAC
CAACCCCCGCCATCCTACCCAGCACACACACACCCTGCTAAC
AAAGACACCCCCCACAGTTATGTAGCTTACCTCCTCAAAGCAAT
NNNNNNNNNNNNNNNNNNNNNNNNNNNNNNNGATCACAGGTCTATC
ACCCTATTAACCACTCACGGGAGCTCTCCA
>scaffold#2
GTATGCACGCGATAAGCATTGCAGACGCTGGAGCCGGAGCACCCT
CTGCCTCATCCTATTATTATCGCACCTACGTTCAATATTACAGG
NNNNNNNNNNNNNNNNNNNNNNNNNNNNNNNNNATTAATTATGCT
GTAGGACATAATAATAACAATTGAATGTCTGATAACAAAAAATTC
CACCAAACCCCCCTCCCCGCTCTGGCCNNNNNACATCACC
CATAAACAAATAGGTTGGTCCTAGCCTTCTATTAGCATCCCCT
TCCAGTGAGTTCACCCCTCTAAATCACCACGATCAAATGCAGCTA
AAACGCTTAGCCTAGCCACACCCCCACGGGAAACACGAAAGTTA
ACTAAGCTATAACTAACCCAGGGT
```

Scaffolds

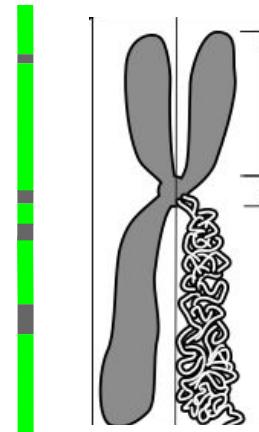


Important definitions

```
>chr22
GTATGCACGCGATAGCATTGCGAGACGCTGGAGCCGGAGCACCCCT
T
CTGCCTCATCCTATTATTATCGCACCTACGTTCAATATTACAGG
G
NNNNNNNNNNNNNNNNNNNNNNNNNNNNNNNNNNTAAATTAATGCT
TGTAGGACATAATAACAATTGAATGTCTGATAACAAAAAATT
TCCACCAAACCCCCCTCCCCGCTCTGGCCNNNNNNACATCA
CCCCATAAACAAATAGTTGGTCCTAGCCTTCTATTAGCATCC
CCGTTCCAGTGAGTTCACCTCTAAATCACCACGATAAAATGCA
GCTAAAACGCTTAGCCTAGCCACACCCCCACGGGAAACACGAAA
GTTTAACTAAGCTATACTAACCCCAGGGTNNNNNNAACCCCCAAA
AACAAAGAACCTAACACCCAGCCTAACCGAGATTCAACAG
TCACCCCCCAACTAACACATTATTTCCCTCCACCAACCCCCG
CCCACCTACCCAGCACACACACCGCTGCTAACCAAAGACACC
CCCCACAGTTATGTAGCTACCTCCTCAAAGCAATNNNNNNNNN
NNNNNNNNNNNNNNNNNNNNNGATCACAGGTCTATCACCTATTA
ACCACTCACGGGAGCTCCA
```

Chromosom

e



Important definitions

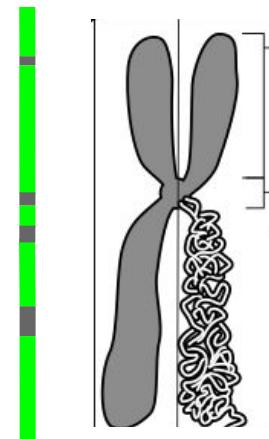
Contigs



Scaffolds



Chromosome



Which approaches?

- Greedy (“Simple” approach)
- Overlap-Layout-Consensus (OLC)
- de Bruijn graphs

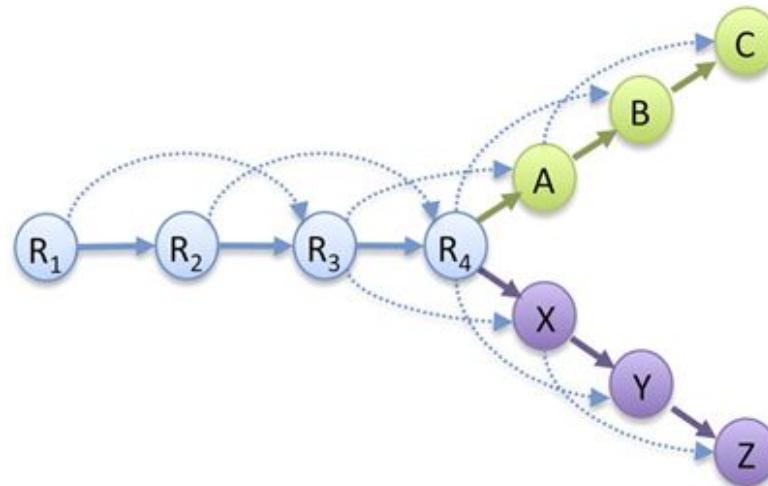
Simple approach - Greedy

- Principle:
 1. Pairwise alignment of all reads
 2. Identify fragments that have largest overlap
 3. Merge these
 4. Repeat until all overlaps are used
- Can only resolve repeats smaller than read length
- High computational cost with increasing no. reads

Overlap-Layout-Consensus

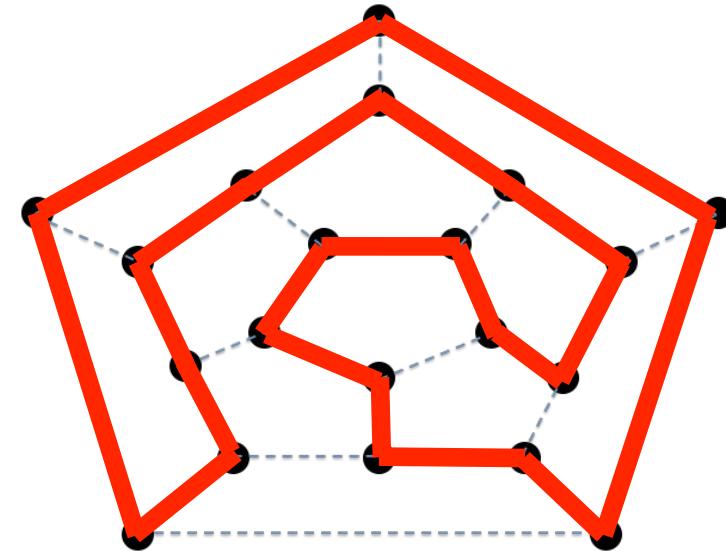
- Create overlap graph by all-vs-all alignment (Overlap)
- Build graph where each node is a read, edges are overlaps between reads (Layout)

| | |
|------------------|----------|
| R ₁ : | GACCTACA |
| R ₂ : | ACCTACAA |
| R ₃ : | CCTACAAG |
| R ₄ : | CTACAAGT |
| A: | TACAAGTT |
| B: | ACAAGTTA |
| C: | CAAGTTAG |
| X: | TACAAGTC |
| Y: | ACAAGTCC |
| Z: | CAAAGTCG |



Overlap-Layout-Consensus

- Create consensus sequence
- We need to use **graph theory** to solve the graph
- Find the *Hamiltonian path*
- i.e. visit each node *exactly once*



Imagine trying to solve this for a graph of hundred of thousands of nodes (=reads)

Overlap-Layout-Consensus

- Not good with many short reads -> lots of alignment!
- With short read lengths, hard to resolve repeats
- Good for large read lengths:
 - PacBio, Oxford Nanopore, 10X Genomics, 454, Ion Torrent, Sanger
- Example assemblers: Canu, Celera, Newbler

de Bruijn graph

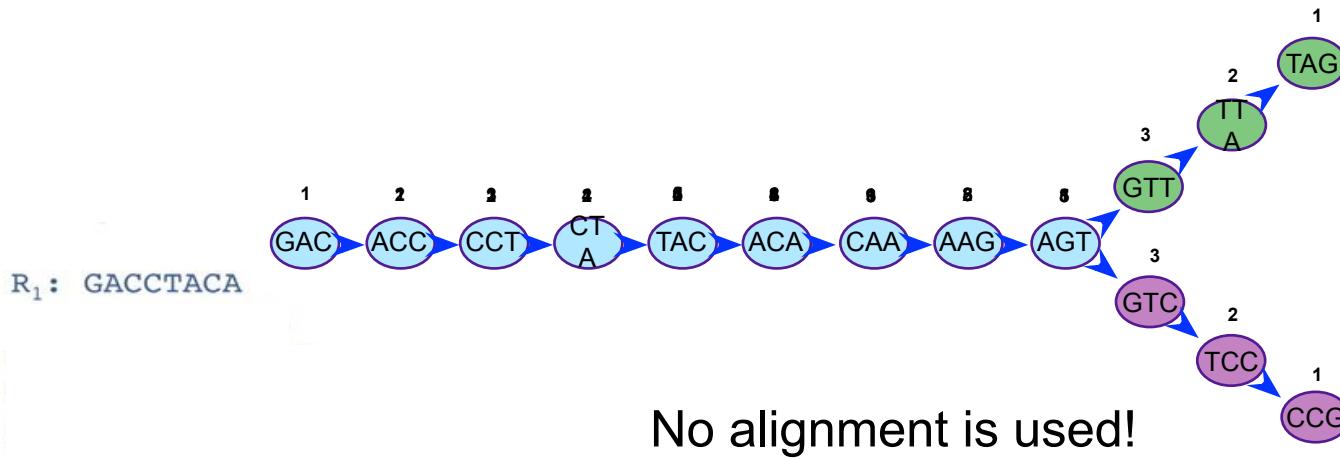
- Directed graph of overlapping items (here DNA sequences)
- Instead of comparing reads, decompose reads into k -mers
 - Graph is created by mapping the k -mers to the graph
 - Each k -mer only exists once in the graph
 - Problem is reduced to walking Eulerian path (visiting each edge once) - this is a solveable problem

Drawbacks ...

- Lots of RAM required (**1-1000 GB !**)
- Optimal k can not be identified *a priori*, must be experimentally tested for each dataset
- small k : very complex graph, large k : limited overlap in low coverage areas
- Iterative approach to find best assembly

How is the graph constructed?

- Same 10 reads, extract k -mers from reads and map onto graph, $k = 3$:



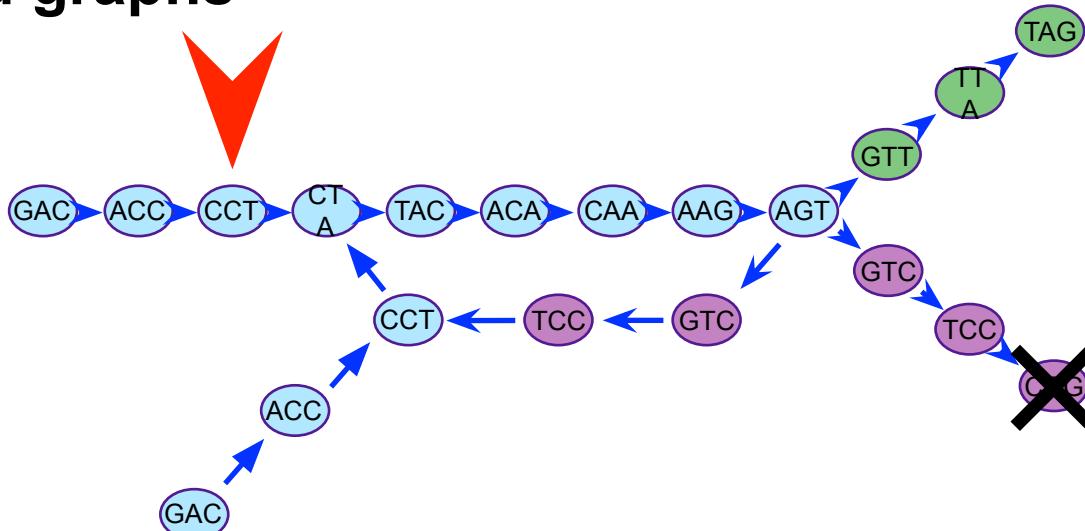
Different assemblers uses
different modifications of the de
Bruijn graphs

Complicated graphs

R₁: GACCTACA
R₂: ACCTACAA
R₃: CCTACAAG
R₄: CTACAAGT

A: TACAAGTT
B: ACAAGTTA
C: CAAGTTAG

X: TACAAGTC
Y: ACAAGTCC
Z: CAAGTCCT

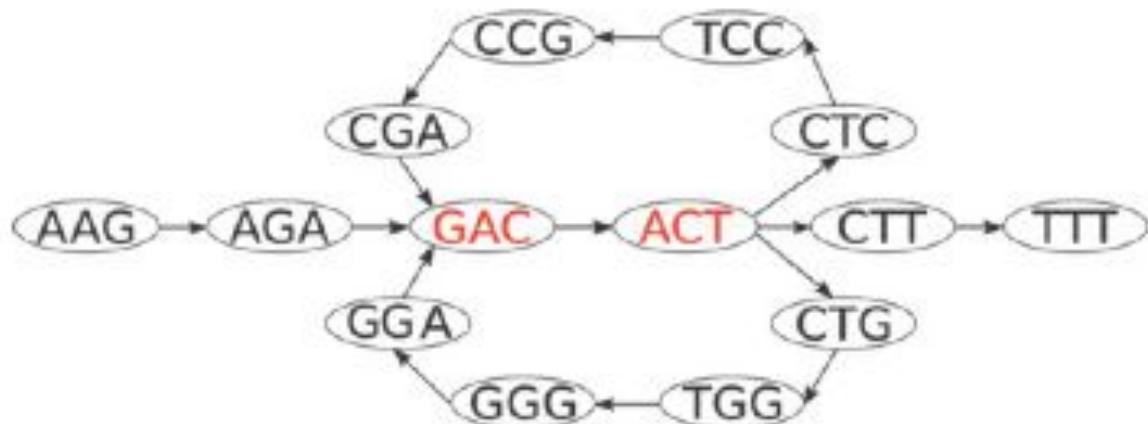


Large genomes with many repeats/errors creates very large graphs

Create the *de Bruijn* graph of this genome using
 $k=3$

AAGACTCCGACTGGGACTTT

AA**GACT**CC**GACT**GGG**GACT**TT

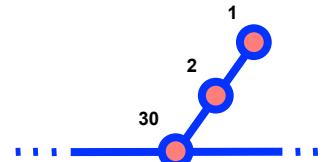


A de Bruijn graph of a sequence

After building: Simplify

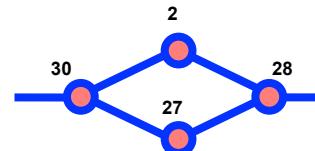
Clip tips

(seq err,end)

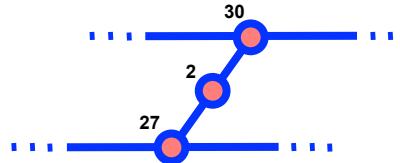


Pinch bubbles

(seq err, middle,
SNP)

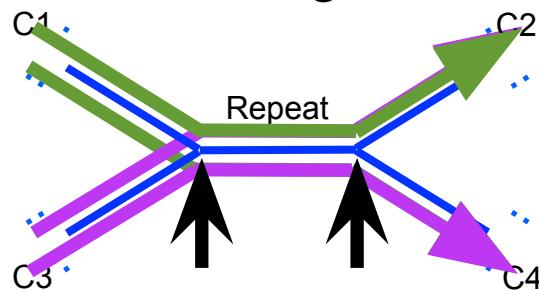


Remove low cov.
links

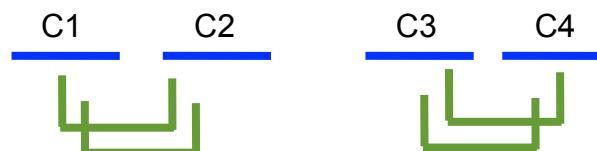


Create contigs and scaffolds

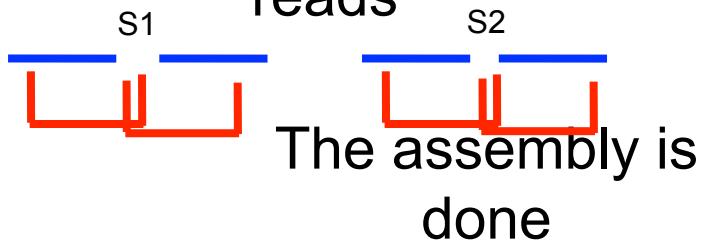
Cut graph at repeat boundaries to create contigs



Use paired-end information to resolve repeats and combine to scaffolds



Fill potential gaps using PE reads

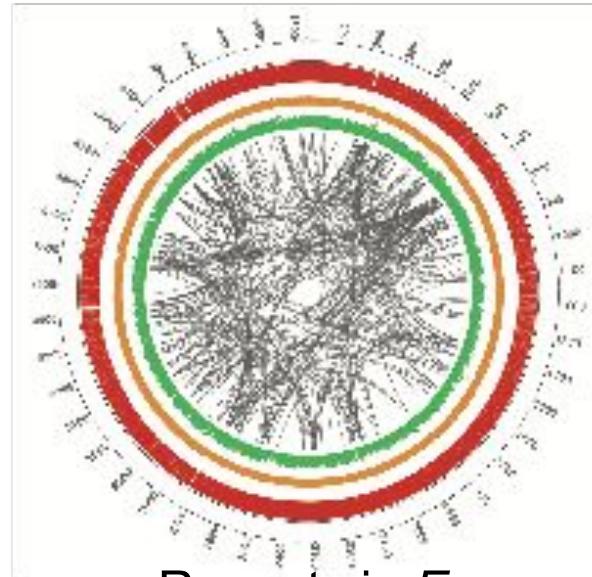


Iterate parameters

- Re-run with different k -sizes, find optimum
- Run with multiple k -mers at the same time! (eg. SPAdes)
- Compare assembly statistics such as, assembly length, N50, no. contigs
- Assembly refinement
 - Break contigs not supported by PE/MP reads
 - Analyze assembly using REAPR or QUAST

Successful *de novo* assembly

- Success is a factor of:
 - Genome size, **genomic repeats(!)**, ploidy
 - High coverage, long read lengths, PE/MP libraries



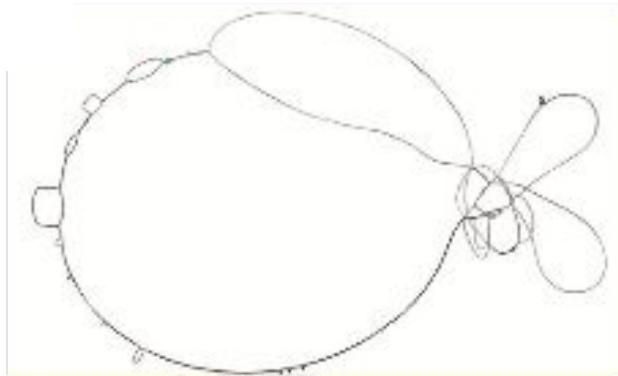
Repeats in *E.
coli*

Improving *de novo* assemblies

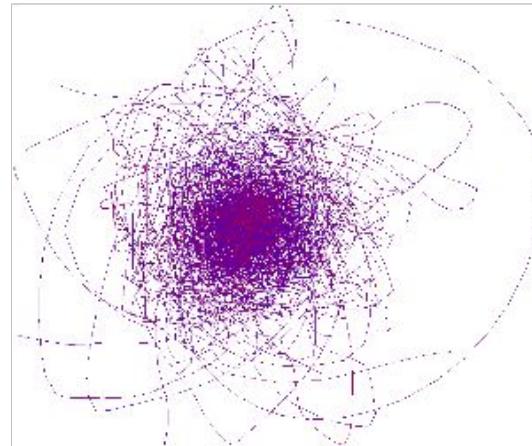
- Paired-end & Mate-pair for long range continuity
- Hybrid approaches (combine Illumina with PacBio/Oxford Nanopore)
- Synthetic long reads: Illumina Synthetic Reads (Moleculo) or 10X Genomics
- Hi-C contact maps

Two bacterial genomes *de Bruijn* graphs

Few
repeats



“more”
repeats



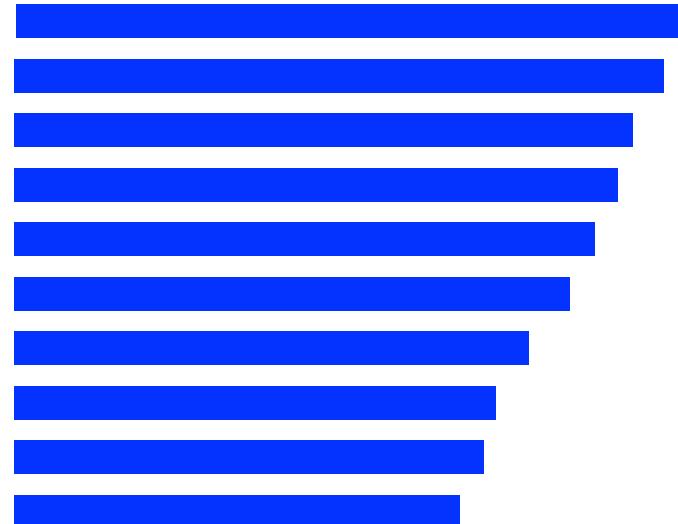
Flicek & Birney, Nat.Methods 2009

Zerbino, 2009

N50: Assembly quality

N50: What is the smallest piece in the largest half of the assembly?

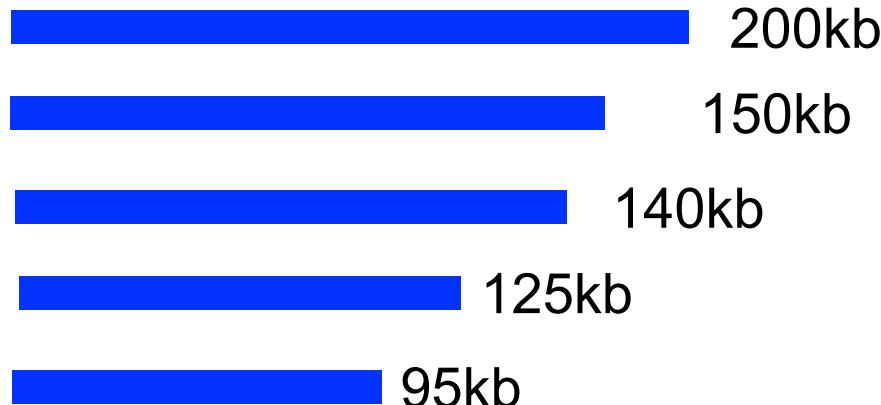
- Calculate sum of assembly
- Order contigs by size
- Sum contigs starting by largest
- When half the sum is reached, N50 is the length of the contig



N50 example

5 scaffolds, calculate

N50:



Sum: $200+150+140+125+95=710\text{kb}$

Half: $710 / 2 = 355\text{kb}$

$$200\text{kb} + 150\text{kb} = 350\text{kb}$$

$$350\text{kb} + 140\text{kb} = 490\text{kb}$$

$490\text{kb} > 355\text{kb} \Rightarrow \mathbf{N50: 140\text{kb}}$

Some assemblers

- OLC: Canu, Newbler
- de Bruijn: Allpaths-LG, SPAdes, Velvet(best), SOAPdenovo, Megahit (very lean), ...
- other: MIRA, SGA, Flye (very good for 3g NGS)

Used in exercises today

Exercise time!